

# Smith Normal Form and Integer solutions to linear equations

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In this tract,  $\text{GCD}(0, \dots, 0)$  is zero; this, since every integer divides zero.

All matrices (and vectors) are **integer-valued**,<sup>♥1</sup> unless specified otherwise. So a square matrix  $\mathbf{R}$  is **invertible**<sup>♥1</sup> IFF  $\text{Det}(\mathbf{R}) \in \{\pm 1\}$ .

For two, say,  $3 \times 5$  matrices  $\mathbf{A}$  and  $\mathbf{B}$ , write:

$$\begin{aligned} \mathbf{A} &\stackrel{r}{\sim} \mathbf{B}, & \text{read "A is row equivalent to B";} \\ \mathbf{A} &\stackrel{c}{\sim} \mathbf{B}, & \text{read "A is column equiv. to B";} \\ \mathbf{A} &\stackrel{rc}{\sim} \mathbf{B}, & \text{read "A is rowcol equivalent to B";} \end{aligned}$$

if there exist *invertible* matrices  $\overset{3 \times 3}{\mathbf{R}}$  and  $\overset{5 \times 5}{\mathbf{C}}$  such that, respectively,

$$\mathbf{RA} = \mathbf{B}; \quad \mathbf{AC} = \mathbf{B}; \quad \mathbf{RAC} = \mathbf{B}.$$

## Smith Normal Form

An  $r \times c$  matrix  $\mathbf{G}$

$$5: \begin{bmatrix} \delta_1 & & & 0 \\ & \delta_2 & & 0 \\ & & \ddots & \vdots \\ & & & \delta_m & 0 \end{bmatrix} \quad \begin{array}{l} \text{(Here, } \mathbf{m} := \text{Min}(\mathbf{r}, \mathbf{c}). \text{ In} \\ \text{this example, } \mathbf{c} \text{ equals } \mathbf{r} + 1. \\ \text{All unshown entries are zero.)} \end{array}$$

is in **Smith form** if its only non-zero entries –the **pivot-values**– are on its main-diagonal. (Zeros are allowed on the main-diagonal). **Smith form** requires that the matrix be  $\mathbb{Z}$ -valued, and that all non-zero values on the diagonal occur *before* the zeros. Let  $\pi = \pi(\mathbf{G})$  denote the number of pivots (non-zero values).

Our  $r \times c$  matrix  $\mathbf{G}$  is in **Smith normal form**

$$6: \begin{bmatrix} \delta_1 & & & 0 \\ & \ddots & & \vdots \\ & & \delta_\pi & \vdots \\ 0 & \dots & 0 & \dots & 0 \end{bmatrix} \quad \begin{array}{l} \text{(In this example, } \mathbf{m} = \mathbf{r} \\ \text{and } \pi = \mathbf{r} - 1. \text{ Thus } \delta_{\mathbf{m}} \text{ is} \\ \text{necessarily 0.)} \end{array}$$

if, in addition, the pivots-values are positive *and*

$$7: \quad \delta_1 \bullet \dots \bullet \delta_{\pi-1} \bullet \delta_\pi \bullet \dots \bullet \delta_{\mathbf{m}}.$$

(This divisibility-condition forces zeros on the diagonal to occur last.)

<sup>♥1</sup>More generally, our matrices' entries come from a *euclidean domain*,  $ED$ . An *invertible*  $\mathbf{R}$  has  $\text{Det}(\mathbf{R}) \in \text{Units}(ED)$ . For row operations, we may: *Add any ED-multiple of a row to another; Multiply a row by any ED-unit.* Ditto for column-ops.

**Elementary row operations.** Applied to an  $r \times c$  matrix  $\Gamma$ , the (elementary) row-operations<sup>♥1</sup> are

- a: Exchanging two rows.
- b: Adding a  $\mathbb{Z}$ -multiple of one row to another.
- c: Multiplying a row by  $-1$ .

Applying a row-op to  $\Gamma$  produces  $\mathbf{E}\Gamma$ , where  $\mathbf{E}$  is an  $r \times r$  **elementary matrix**. Notice that  $\text{Det}(\mathbf{E})$  is  $\pm 1$ , since we are allowed to multiply a row only<sup>♥1</sup> by  $-1$ .

Analogously, there are the **column operations**. Applying a col-op to  $\Gamma$  produces  $\Gamma\hat{\mathbf{E}}$ , where  $\hat{\mathbf{E}}$  is an  $c \times c$  elementary matrix.

Applying  $j$ -many row-ops and  $k$ -many col-ops to  $\Gamma$  produces a matrix

$$\mathbf{G} := \mathbf{R}\Gamma\mathbf{C}, \quad \begin{array}{l} \text{where } \mathbf{R} := \mathbf{E}_j \dots \mathbf{E}_2 \mathbf{E}_1 \\ \text{and } \mathbf{C} := \hat{\mathbf{E}}_1 \hat{\mathbf{E}}_2 \dots \hat{\mathbf{E}}_k. \end{array}$$

This  $\mathbf{R}$  is an integer matrix with  $\text{Det}(\mathbf{R}) \in \{\pm 1\}$ . Ditto  $\mathbf{C}$ . These are the *bookkeeping* matrices; our  $\mathbf{R}$  keeps track of the (cumulative) row-ops, and  $\mathbf{C}$  keeps track of the col-ops. Converting  $\Gamma$  to  $\mathbf{G}$  via elem. row-ops and col-ops manifests *rowcol-equivalence*.

It turns out that each matrix  $\Gamma$  is rowcol-equivalent to a *unique* Smith-Normal-Form matrix. We write this as  $\mathbf{G} := \text{SNF}(\Gamma)$ .

**8: Smith Normal Form Thm.** *Each matrix  $\Gamma$  is rowcol-equiv to some SNF-matrix,  $\mathbf{G}$ . Moreover, the SNF is unique. (I.e, no two distinct SNFs are rowcol-equiv.)*  $\diamond$

**Existence of SNF.** Applying Lemma 10, below, to our  $r \times c$  matrix  $\Gamma$ , yields an integer  $\alpha_1$  and matrix  $\mathbf{B}$  with

$$\dagger: \quad \Gamma \stackrel{rc}{\sim} \begin{bmatrix} \alpha_1 & 0 & 0 & \dots & 0 \\ 0 & & & & \\ \vdots & & \mathbf{B} & & \\ 0 & & & & \end{bmatrix}.$$

By induction on the dimensions of a matrix, our

$$\ddagger: \quad \mathbf{B} \stackrel{rc}{\sim} \begin{bmatrix} \alpha_2 & & \\ & \ddots & \\ & & \alpha_{\mathbf{m}} \end{bmatrix}, \quad \text{where } \mathbf{m} := \text{Min}(\mathbf{r}, \mathbf{c}).$$

Finally, the rowcol equivalence of ( $\ddagger$ ) can be done *inside* of ( $\dagger$ ), thanks to the zeros in the first row and first column of  $\text{RhS}(\ddagger)$ . Consequently

$$\Gamma \stackrel{rc}{\sim} \begin{bmatrix} \alpha_1 & \alpha_2 & & \\ & \ddots & & \\ & & & \alpha_{\mathbf{m}} \end{bmatrix}.$$

Our last obligation to to arrange divisibility. Apply **Lemma 9** to the 1:2 submatrix (meaning, the  $2 \times 2$  matrix formed by rows and columns 1 and 2) of  $\begin{bmatrix} \alpha_1 & \cdots & \alpha_m \\ \vdots & \ddots & \vdots \\ \alpha_1 & \cdots & \alpha_m \end{bmatrix}$ . Now apply **Lemma 9** to 1:3, then 1:4, and continue up to the 1: $m$  submatrix. Letting  $g := \text{GCD}(\alpha_1, \dots, \alpha_m)$ , we have shown that

$$\Gamma \stackrel{rc}{\sim} \begin{bmatrix} g & & \\ & \beta_2 & \\ & & \ddots \\ & & & \beta_m \end{bmatrix},$$

where each  $\beta$  number is an integer multiple of  $g$ . Now proceed inductively on the  $\beta$ -submatrix.  $\blacklozenge$

**9: Lemma.** Suppose  $\alpha, \beta \in \mathbb{Z}$ . Then

$$\begin{bmatrix} \alpha & 0 \\ 0 & \beta \end{bmatrix} \stackrel{rc}{\sim} \begin{bmatrix} g & 0 \\ 0 & \ell \end{bmatrix},$$

where  $g := \text{GCD}(\alpha, \beta)$  and  $\ell := \text{LCM}(\alpha, \beta)$ .  $\blacklozenge$

**Proof.** Since  $\mathbb{Z}$  is a Euclidean domain, there exist integers  $S, T$  st.  $S\alpha + T\beta = g$ . Thus

$$\begin{bmatrix} \alpha & 0 \\ 0 & \beta \end{bmatrix} \stackrel{c}{\sim} \begin{bmatrix} \alpha & S\alpha \\ 0 & \beta \end{bmatrix} \stackrel{r}{\sim} \begin{bmatrix} \alpha & S\alpha + T\beta \\ 0 & \beta \end{bmatrix} \stackrel{c}{\sim} \begin{bmatrix} g & \alpha \\ \beta & 0 \end{bmatrix}.$$

Take integers  $J, K$  with  $\alpha = -Jg$  and  $\beta = Kg$ . Then

$$\begin{bmatrix} g & \alpha \\ \beta & 0 \end{bmatrix} \stackrel{c}{\sim} \begin{bmatrix} g & \alpha + Jg \\ \beta & J\beta \end{bmatrix} = \begin{bmatrix} g & 0 \\ \beta & J\beta \end{bmatrix} \stackrel{r}{\sim} \begin{bmatrix} g & 0 \\ \beta - Kg & J\beta \end{bmatrix} = \begin{bmatrix} g & 0 \\ 0 & JKg \end{bmatrix}.$$

And  $-JKg = \frac{-Jg \cdot Kg}{g} = \frac{\alpha\beta}{g} = \pm\ell$ . So negate column-two if need be.  $\blacklozenge$

**10: Lemma.** Suppose  $A$  is an  $\mathbf{r} \times \mathbf{c}$  matrix with  $\mathbf{r}, \mathbf{c} \geq 1$ . Then

$$A \stackrel{rc}{\sim} \begin{bmatrix} \alpha & 0 & \cdots & 0 \\ \vdots & & & \\ 0 & & & B \end{bmatrix}$$

for some integer  $\alpha$  and  $[\mathbf{r}-1] \times [\mathbf{c}-1]$  matrix  $B$ .  $\blacklozenge$

**Proof.** Write  $A = [a_{ij}]_{i,j}$ . Suppose we have a position  $\sigma = (i_0, j_0)$  where  $a_\sigma \neq 0$ . For a different position on the same column,  $\tau = (i, j_0)$  with  $i \neq i_0$ , divide  $a_\sigma$  into  $a_\tau$  to get a quotient and remainder (integers):

$$a_\sigma = Q \cdot a_\tau + \mathcal{R}, \quad \text{with } |\mathcal{R}| < |a_\sigma|.$$

Subtracting  $Q \cdot \text{Row}(i_0, A)$  from  $\text{Row}(i, A)$  yields a new matrix  $A' \stackrel{r}{\sim} A$  with  $|a'_\tau| < |a_\sigma|$ .

For each psn  $\tau \in \text{ColOrRow}(\sigma)$  with  $\tau \neq \sigma$ , we can subtract a multiple of  $\sigma$ 's row or column, obtaining a

11: New matrix  $A' \stackrel{rc}{\sim} A$  for which  $\forall \tau \in \text{ColOrRow}(\sigma)$  with  $\tau \neq \sigma$ :  $|a'_\tau| < |a'_\sigma|$ .

**Iterating.** Let  $\text{Min}(A)$  be the minimum of  $|a_\sigma|$ , taken over all positions  $\sigma$  for which  $a_\sigma \neq 0$ .

To prove the lemma, WLOG  $\text{Min}(A) \neq 0$ . Apply the following procedure.

i: Set  $\mu := \text{Min}(A)$  and pick  $\sigma$  with  $|a_\sigma| = \mu$ .

ii: Apply (11) to produce a matrix  $A' \stackrel{rc}{\sim} A$  with  $0 < \text{Min}(A') \leq \mu$ .

iii: If  $\text{Min}(A') = \mu$ , then **Stop!** Else, set  $A := A'$ , and go to step (i).

We will eventually **Stop**, since the  $\mu$ -numbers form a decreasing sequence of posints. When we stop, we have, for the current position  $\sigma$ , that

$$\text{Min}(A') = \text{Min}(A) = |a_\sigma|.$$

But operation (11) forces *all* the other entries in  $\sigma$ 's row and col to have smaller absolute value than  $a_\sigma$ . So the only way that (11) could have *failed* to lower the matrix's  $\text{Min}()$ , is if

$$\forall \tau \in \text{ColOrRow}(\sigma) \text{ with } \tau \neq \sigma: \quad a'_\tau = 0.$$

And this says that  $\sigma$  is a *pivot-position* for  $A$ . Simply swap rows and swap columns so that the pivot is now in position  $(1, 1)$ . Voila RhS(10).  $\blacklozenge$

**Setting-up uniqueness of SNF.** Imagine an  $\mathbf{r} \times \mathbf{c}$  matrix  $A$ . There are  $\binom{\mathbf{r}}{3}$  many<sup>2</sup> tripletons  $\mathcal{R} \subset [1.. \mathbf{r}]$ , and  $\binom{\mathbf{c}}{3}$  many tripletons  $\mathcal{C} \subset [1.. \mathbf{c}]$ . Let  $A_{\mathcal{R} \times \mathcal{C}}$  denote the induced  $3 \times 3$  submatrix. Thus

$$S_3^A := \left\{ \text{Det}(A_{\mathcal{R} \times \mathcal{C}}) \mid \begin{array}{l} \mathcal{R} \subset [1.. \mathbf{r}] \text{ is a tripleton,} \\ \text{and so is } \mathcal{C} \subset [1.. \mathbf{c}] \end{array} \right\}$$

<sup>2</sup>This  $\binom{\mathbf{r}}{3}$  is the *binomial coefficient* “ $\mathbf{r}$  choose 3”. It is the number of ways of choosing 3 objects from  $\mathbf{r}$ -many distinct objects. E.g,  $\binom{5}{3}$  equals 10.

is a set of integers. Our goal is to show that

$$\mathcal{G}_3^A := \text{GCD}(S_3^A)$$

is an *invariant* of row-column equivalence.

**12: Invariance Lemma.** Consider matrices  $A \stackrel{rc}{\sim} B$ .  
Then  $\mathcal{G}_k^A = \mathcal{G}_k^B$ , for each  $k = 1, 2, 3, \dots$   $\diamond$

**Pf of uniqueness.** Courtesy (13†), sequence  $(\mathcal{G}_k^A)_{k=1}^m$  determines the  $(\delta_k)_{k=1}^m$  sequence.  $\blacklozenge$

**Commentary.** Although  $\mathbf{G} := \text{SNF}(\mathbf{A})$  is unique, there can be *many* pairs  $(\mathbf{R}, \mathbf{C})$  such that  $\mathbf{RAC}$  equals  $\mathbf{G}$ . Extreme cases are  $\mathbf{A} := \mathbf{I}$ , or  $\mathbf{A}$  is the zero-matrix.  $\square$

**Pf.** WELOG  $k = 3$ . WELOG, we obtained  $\mathbf{B}$  from  $\mathbf{A}$

by adding  $98 \cdot \text{Row}(\mathbf{A}, 7)$  to  $\text{Row}(\mathbf{A}, 5)$ .

So ISTShow that  $\boxed{\mathcal{G}_3^A \mid \mathcal{G}_3^B}$ ; for by symmetry, then,  $\mathcal{G}_3^A \mid \mathcal{G}_3^B$ . BTWay, we've done a row-op, so the set  $\mathcal{C}$  plays no essential role. Hence, let  $\mathbf{A}_{\mathcal{R}}$  and  $\mathbf{B}_{\mathcal{R}}$  mean  $\mathbf{A}_{\mathcal{R} \times \mathcal{C}}$  and  $\mathbf{B}_{\mathcal{R} \times \mathcal{C}}$ , for some particular choice of  $\mathcal{C}$ .

Consider a tripleton  $\mathcal{R} \subset [1..r]$ . If  $\mathcal{R} \not\ni 5$ , then  $\mathbf{B}_{\mathcal{R}} = \mathbf{A}_{\mathcal{R}}$ .

In contrast, suppose  $\boxed{\mathcal{R} \ni 5}$ . If  $\mathcal{R} \ni 7$ , then

$$\text{Det}(\mathbf{B}_{\mathcal{R}}) = \text{Det}(\mathbf{A}_{\mathcal{R}}).$$

Now suppose  $\mathcal{R} \not\ni 7$ . Writing  $\mathcal{R}$  as  $\{5, i, i'\}$ , then,

$$\text{Det}(\mathbf{B}_{\mathcal{R}}) = \text{Det}(\mathbf{A}_{\mathcal{R}}) + 98 \cdot \text{Det}(\mathbf{A}_{\{7, i, i'\}}).$$

And  $\mathcal{G}_3^A$  divides both  $\text{Det}(\mathbf{A}_{\mathcal{R}})$  and  $\text{Det}(\mathbf{A}_{\{7, i, i'\}})$ .

In all three cases, we have that  $\mathcal{G}_3^A \mid \text{Det}(\mathbf{B}_{\mathcal{R}})$ .  $\blacklozenge$

**13: Coro.: Uniqueness of SNF.** For an  $r \times c$  matrix  $\mathbf{A}$ , let  $\mathbf{m} := \text{Min}(\mathbf{r}, \mathbf{c})$ . Let  $\mathbf{G}$  be some SNF of  $\mathbf{A}$ , as in (6). And let  $\pi := \text{Rank}(\mathbf{A})$ , i.e, the number of pivots in  $\mathbf{G}$ . Then  $\mathbf{A}$  has only one SNF, since

For each  $k = 1, 2, \dots, \mathbf{m}$ :

13†:  $\mathcal{G}_k^A \stackrel{\text{note}}{=} \mathcal{G}_k^G$  equals the product  $\delta_1 \cdot \delta_2 \cdots \delta_k$ .

In particular  $\mathcal{G}_1^A \mid \dots \mid \mathcal{G}_m^A$ . Indeed, the ratios

$$13‡: \frac{\mathcal{G}_2^A}{\mathcal{G}_1^A} \mid \frac{\mathcal{G}_3^A}{\mathcal{G}_2^A} \mid \dots \mid \frac{\mathcal{G}_\pi^A}{\mathcal{G}_{\pi-1}^A}. \quad \blacklozenge$$

**Pf of (13†).** Our  $\mathbf{G}$  is as in (6) and (7). To compute, say,  $\mathcal{G}_3^G$ , take a tripleton  $\mathcal{R} = \{i, i', i''\}$  with  $i < i' < i''$ . For  $\text{Det}(\mathbf{G}_{\mathcal{R} \times \mathcal{C}})$  to be non-zero, necessarily  $\mathcal{C} = \mathcal{R}$ . Thus  $\text{Det}(\mathbf{G}_{\mathcal{R} \times \mathcal{C}})$  equals  $\delta_i \cdot \delta_{i'} \cdot \delta_{i''}$ . Now use (7).  $\blacklozenge$

## A system of linear equations over $\mathbb{Z}$

(Matrices are required to be integer-valued.) Consider a matrix-equation of the form

$$14: \quad \overset{r \times c}{\Gamma} \cdot \overset{c \times 1}{\Sigma} = \overset{r \times 1}{\Upsilon}.$$

The *coefficient matrix* is  $\Gamma$ , and  $\Upsilon$  is a *target matrix*. We want to describe the set (it will turn out to be a  $\mathbb{Z}$ -lattice) of target matrices. And given a target, we seek the set of *solution matrices*  $\Sigma$  for (14).

Here is the logic. There may be many solutions,  $S$ , to  $\Gamma S = \Upsilon$ , one of which is  $\Sigma$ . But this *augmented* system

$$\begin{aligned} \Gamma \cdot S &= \Upsilon & \text{and} \\ \mathbf{I} \cdot S &= \Sigma \end{aligned}$$

(where  $\mathbf{I}$  represents an identity matrix) has the *unique* solution  $S := \Sigma$ . We solve for  $S$  using row and column operations. We'll then discover "free variables" which allow us to describe *all* solns  $S$  to  $\Gamma S = \Upsilon$ .

**Sequences of matrices.** Initialize matrices

$$\begin{aligned} R_0 &:= \overset{r \times r}{\mathbf{I}}, & G_0 &:= \overset{r \times c}{\Gamma}, & C_0 &:= \overset{c \times c}{\mathbf{I}} \\ \text{and column-vectors } S_0 &:= \overset{c \times 1}{\Sigma}, & U_0 &:= \overset{r \times 1}{\Upsilon}, \end{aligned}$$

Row&column operations will modify these, always retaining that

$$\begin{aligned} 1_n, 2_n: & \quad G_n S_n = U_n & \text{and} & \quad U_n = R_n \Upsilon. \\ 3_n, 4_n: & \quad C_n S_n = \Sigma & \text{and} & \quad G_n = R_n \Gamma C_n. \end{aligned}$$

The Row&col ops will be effectuated by elementary matrices  $E$  with integer entries and  $\text{Det}(E) \in \{\pm 1\}$ , i.e in  $\text{Units}(\mathbb{Z})$ .

**Row operations.** A row-op matrix  $\overset{r \times r}{E}$  updates the matrix-sequences as follows:

$$\begin{aligned} 1, 2: & \quad \overset{G_{n+1}}{\underbrace{E G_n}} S_n = \overset{U_{n+1}}{\underbrace{E U_n}} & \text{and} & \quad \overset{U_{n+1}}{\underbrace{E U_n}} = \overset{R_{n+1}}{\underbrace{E R_n}} \Upsilon. \\ 4: & & \text{and} & \quad \overset{E G_n}{\underbrace{\quad}_{G_{n+1}}} = \overset{E R_n}{\underbrace{\quad}_{R_{n+1}}} \Gamma C_n. \end{aligned}$$

The other matrices retain their value, i.e  $C_{n+1} := C_n$ .

**Column ops.** A column-op  $\overset{c \times c}{E}$  updates like this:

$$\begin{aligned} 1: & \quad \overset{G_{n+1}}{\underbrace{G_n E}} \overset{S_{n+1}}{\underbrace{E^{-1} S_n}} = U_n & \text{and} \\ 3, 4: & \quad \overset{C_n E}{\underbrace{\quad}_{C_{n+1}}} \overset{S_n}{\underbrace{E^{-1} S_n}_{S_{n+1}}} = \Sigma & \text{and} & \quad \overset{G_n E}{\underbrace{\quad}_{G_{n+1}}} = R_n \Gamma \overset{C_n E}{\underbrace{\quad}_{C_{n+1}}}. \end{aligned}$$

**Result.** We rowcol-reduce  $G_0 = \Gamma$  to a Smith Form. Supposing this happens at  $n=57$ , I use boldface letters to denote

$$\begin{aligned} \mathbf{R} &:= R_{57}, & \mathbf{G} &:= G_{57}, & \mathbf{C} &:= C_{57} \\ \mathbf{S} &:= S_{57}, & \mathbf{U} &:= U_{57}. \end{aligned}$$

I henceforth use " $\infty$ " to denote this last stage "57".

Since we preserved  $([1, 2, 3, 4]_n)$ , we now have

$$\begin{aligned} 1_{\infty}, 2_{\infty}: & \quad \mathbf{G} \mathbf{S} = \mathbf{U} & \text{and} & \quad \mathbf{U} = \mathbf{R} \Upsilon. \\ 3_{\infty}, 4_{\infty}: & \quad \mathbf{C} \mathbf{S} = \Sigma & \text{and} & \quad \mathbf{G} = \mathbf{R} \Gamma \mathbf{C}. \end{aligned}$$

The improvement over  $([1, 2, 3, 4]_0)$  is that, now, the (new) coefficient matrix  $\mathbf{G}$  is in Smith-form, e.g (6).

**Pivot- and free-columns.** Let  $\pi$  be the number of pivots in  $\mathbf{G}$ ; note that  $\pi = \text{Rank}(\Gamma)$ . Use  $\varphi := c - \pi$  for the number of *free columns*.

Decompose  $\mathbf{G}$  into its lefthand  $\pi$ -many columns, call it  $\overset{r \times \pi}{\mathbf{D}}$ , and its righthand  $\varphi$ -many columns:

$$15: \quad \mathbf{G} = \left[ \begin{array}{c|c} \overset{\mathbf{D}}{\begin{matrix} \delta_1 & & \\ & \ddots & \\ & & \delta_\pi \end{matrix}} & \overset{\text{All zero}}{\begin{matrix} 0 & 0 \\ \vdots & \vdots \\ 0 & 0 \end{matrix}} \right]$$

$$16: \quad \mathbf{C} = \left[ \begin{array}{c|c} \mathbf{P} & \mathbf{F} \end{array} \right]$$

Split  $\mathbf{C}$  into its *pivot part*  $\overset{c \times \pi}{\mathbf{P}}$ , and its *free part*  $\overset{c \times \varphi}{\mathbf{F}}$ .

**Divisibility.** From  $([1, 2]_{\infty})$  and (15), an integer column-vector  $\Upsilon$  admits a solution  $\Sigma$  **iff** for each  $i = 1, 2, \dots, \mathbf{m}$ :  $\text{Row}(i, \mathbf{R}) \cdot \Upsilon$  is divisible by  $\delta_i$ .

Another way to state this is:

$$\begin{aligned} 17: & \quad \forall i \in [1 \dots \pi] : \text{Row}(i, \mathbf{R}) \cdot \Upsilon \mid \delta_i, & \text{and} \\ & \quad \forall i \in (\pi \dots \mathbf{r}] : \text{Row}(i, \mathbf{R}) \cdot \Upsilon = 0. \end{aligned}$$

## Lattices

These matrices will give us explicit  $\mathbb{Z}$ -bases for the target lattice and the nullspace lattice.

**18: Nullspace Theorem.** *The  $\mathbb{Z}$ -lattice of vectors  $\Sigma$  such that  $\Gamma\Sigma = \mathbf{0}_{\mathbf{r} \times 1}$ , is  $\varphi$ -dimensional. And the set of columns of  $\mathbf{F}$  is a  $\mathbb{Z}$ -basis for the nullspace-lattice.  $\diamond$*

**Proof.** Set  $\Upsilon := \mathbf{0}_{\mathbf{r} \times 1}$  in  $([1, 2, 3, 4]_\infty)$ . Thus  $\mathbf{U} = \mathbf{0}_{\mathbf{r} \times 1}$ . So  $\mathbf{S}$  solves  $(1_\infty)$  exactly when the top  $\pi$ -many entries in  $\mathbf{S}$  are zero; the rest can vary freely. The resulting set of products  $\Sigma = \mathbf{C}\mathbf{S}$  is exactly the set of products

$$\mathbf{F} \cdot \begin{bmatrix} s_{\pi+1} \\ \vdots \\ s_{\mathbf{c}} \end{bmatrix}, \quad \begin{array}{l} \text{as the } \varphi\text{-many entries} \\ s_{\pi+1}, s_{\pi+2}, \dots, s_{\mathbf{c}-1}, s_{\mathbf{c}} \\ \text{vary freely.} \end{array}$$

Finally, the  $\mathbf{F}$ -columns are linearly-independent, since those of  $\mathbf{C}$  are, since those of  $\mathbf{C}_0$  were.  $\diamond$

**A  $\mathbb{Z}$ -basis for the set of targets.** (Below, use  $\mathbf{M}^t$  to indicate  $\mathbf{M}$ -transpose.) We will show that

$$19: \quad \Gamma\mathbf{P} = \mathbf{R}^{-1}\mathbf{D}, \quad \text{i.e.} \quad \begin{matrix} \mathbf{r} \times \mathbf{r} & \mathbf{r} \times \mathbf{c} & \mathbf{c} \times \pi & & \mathbf{r} \times \pi \\ \mathbf{R} & \Gamma & \mathbf{P} & = & \mathbf{D} \end{matrix},$$

on our way to establishing our main goal.

**20: Target Thm.** *The set of  $\Upsilon$  for which there exists  $\Sigma$  with  $\Gamma\Sigma = \Upsilon$ , is a  $\pi$ -dim'al  $\mathbb{Z}$ -lattice. A  $\mathbb{Z}$ -basis for this lattice is the set of columns of  $\Gamma\mathbf{P}$ , i.e., of  $\mathbf{R}^{-1}\mathbf{D}$ .  $\diamond$*

**Proof of (20) and (19).** Recall that  $\Upsilon = \Gamma\Sigma = \Gamma\mathbf{C}\mathbf{S}$ . We get *each* target as a product

$$\dagger: \quad \Gamma\mathbf{C} \cdot [s_1 \dots s_\pi \ 0 \dots 0]^t,$$

for one particular choice of integer-tuple  $s_1, \dots, s_\pi$ . Why? —because the last  $\varphi$ -many components of  $\mathbf{S}$  are mapped by  $\mathbf{C}$ , bijectively, to the nullspace of  $\Gamma$ .

Courtesy (16), this set of products ( $\dagger$ ) is the set

$$\ddagger: \quad \Gamma\mathbf{P} \cdot [s_1 \dots s_\pi]^t.$$

So we've established (20) —*except* that we still need to prove that  $\mathbf{R}^{-1}\mathbf{D}$  equals  $\Gamma\mathbf{P}$ .

Each  $\mathbf{U}$  vector arises as a product

$$\mathbf{G} \cdot [s_1 \dots s_\pi \ 0 \dots 0]^t,$$

thanks to  $(1_\infty)$  and (15). And this product equals

$$\mathbf{D} \cdot [s_1 \dots s_\pi]^t.$$

Hence  $\Upsilon = \mathbf{R}^{-1}\mathbf{U} = \mathbf{R}^{-1}\mathbf{D} [s_1 \dots s_\pi]^t$ , and this gives the same mapping  $(s_1, \dots, s_\pi) \mapsto \Upsilon$  that  $\Gamma\mathbf{P}$  does. Consequently,  $\mathbf{R}^{-1}\mathbf{D}$  minus  $\Gamma\mathbf{P}$  is the zero-operator —and so they are equal.  $\diamond$

**Mapping targets to solns.** So the  $\mathbf{r} \times \pi$  matrix

$$21: \quad \mathbf{T} := \Gamma\mathbf{P} \stackrel{\text{fact}}{=} \mathbf{R}^{-1}\mathbf{D}$$

maps “*s-tuples*”  $[s_1 \dots s_\pi]^t$  to targets, bijectively, via lefthand multiplication. Thus, in the space of *rational* matrices, our  $\mathbf{T}$  has a *lefthand* inverse.  $\heartsuit^3$

We pick a particular LH-inverse  $\mathbf{T}^\bullet := \mathbf{D}^\bullet \cdot \mathbf{R}$  by specifying the following rational lefthand-inverse to  $\mathbf{D}$ :

$$22: \quad \mathbf{D}^\bullet := \begin{bmatrix} 1/\delta_1 & & 0 \\ & \ddots & \vdots \\ & & 1/\delta_\pi & 0 \end{bmatrix} \in \text{MAT}_{\pi \times \mathbf{r}}(\mathbb{Q}).$$

It satisfies  $\heartsuit^4$  that  $\mathbf{D}^\bullet \mathbf{D} = \mathbf{I}^{\pi \times \pi}$ .

In consequence, lefthand-multiplication by *rational*  $\mathbf{c} \times \mathbf{r}$  matrix

$$23: \quad \mathbf{M} := \mathbf{P}\mathbf{D}^\bullet\mathbf{R}$$

carries each target  $\Upsilon$  to a particular solution  $\Sigma$  for which  $\Gamma\Sigma = \Upsilon$ . For each target  $\Upsilon$ , then,

$$24: \quad \left\{ \Sigma \mid \Gamma\Sigma = \Upsilon \right\} = \mathbf{M}\Upsilon + \mathbf{F} \begin{bmatrix} s_{\pi+1} \\ s_{\pi+2} \\ \vdots \\ s_{\mathbf{c}} \end{bmatrix},$$

as  $s_{\pi+1}, \dots, s_{\mathbf{c}-1}, s_{\mathbf{c}}$  vary over the integers.

We call  $\mathbf{M}$  our *selector matrix*. It is determined by our particular reduction of  $\Gamma$  to Smith Form.

If we reduced  $\Gamma$  to Smith *Normal* form, then each quotient  $\varepsilon_i := \delta_\pi / \delta_i$  is a posint. So

$$25: \quad \begin{aligned} \mathbf{M} &= \frac{1}{\delta_\pi} \cdot \widehat{\mathbf{M}}, \quad \text{where the product} \\ \widehat{\mathbf{M}} &:= \mathbf{P} \begin{bmatrix} \varepsilon_1 & & 0 \dots 0 \\ & \varepsilon_2 & 0 \dots 0 \\ & & \ddots & \vdots \\ & & & \varepsilon_\pi & 0 \dots 0 \end{bmatrix} \mathbf{R} \end{aligned}$$

$\heartsuit^3$  Its set of rational LH-inverses is an  $[\mathbf{r}-\pi]$ -dimensional “flat”, i.e., an affine subspace.

$\heartsuit^4$  Reverse-product  $\mathbf{D}\mathbf{D}^\bullet$  will not equal  $\mathbf{I}^{\mathbf{c} \times \mathbf{c}}$ , unless  $\pi \stackrel{\text{happens}}{=} \mathbf{c}$ .

is an *integer* matrix.

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