

## Primer on Cardinalities

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**1a: Defn: Exponentiation.** Given sets  $D$  and  $C$ , logicians define the symbol  $C^D$  as

$$C^D := \left\{ \text{The set of } \underline{\text{all}} \text{ fncs from } D \rightarrow C \right\}.$$

These functions have **Domain**  $D$  and **Co-domain**  $C$ . Note well that  $C^D$  is a *set of functions*. For example, let  $\Omega := \{\alpha, \beta, \gamma\}$  and  $B := \{0, 1\}$ . Then  $B^\Omega$  comprises the  $2^3 = 8$  many fncs

$$\begin{array}{ll} \alpha, \beta, \gamma \mapsto 0, 0, 0; & \alpha, \beta, \gamma \mapsto 1, 0, 0; \\ \alpha, \beta, \gamma \mapsto 0, 0, 1; & \alpha, \beta, \gamma \mapsto 1, 0, 1; \\ \alpha, \beta, \gamma \mapsto 0, 1, 0; & \alpha, \beta, \gamma \mapsto 1, 1, 0; \\ \alpha, \beta, \gamma \mapsto 0, 1, 1; & \alpha, \beta, \gamma \mapsto 1, 1, 1. \end{array}$$

In contrast,  $\Omega^B$  comprises these  $3^2 = 9$  fncs:

$$\begin{array}{lll} 0, 1 \mapsto \alpha, \alpha; & 0, 1 \mapsto \beta, \alpha; & 0, 1 \mapsto \gamma, \alpha; \\ 0, 1 \mapsto \alpha, \beta; & 0, 1 \mapsto \beta, \beta; & 0, 1 \mapsto \gamma, \beta; \\ 0, 1 \mapsto \alpha, \gamma; & 0, 1 \mapsto \beta, \gamma; & 0, 1 \mapsto \gamma, \gamma. \end{array}$$

Note, for finite sets  $P$  and  $Q$ , that  $|P^Q| = |P|^{|Q|}$ . It is for that reason that logicians use this  $\text{Set}^{\text{Set}}$  notation. **[N.B:** Consider sets  $A \asymp B$  with  $A \neq B$ . Although sets  $A^B$  and  $B^A$  have the same *cardinality*, they are *not* the same set; this, since the fncs in  $A^B$  have  $B$  as their domain, whereas those in  $B^A$  have  $A$  as their domain, yet  $B \neq A$ .]  $\square$

**1b: Defn: Powerset.** The **powerset** of a set  $\Omega$ , written  $\mathcal{P}(\Omega)$ , is the set of *all* subsets of  $\Omega$ . Why do logicians sometimes write  $\{0, 1\}^\Omega$  to mean  $\mathcal{P}(\Omega)$ ?

Well, there is a natural bijection between the two: A function  $f: \Omega \rightarrow \{0, 1\}$  yields a subset  $S_f \subset \Omega$  by  $S_f := \{x \in \Omega \mid f(x) = 1\}$ . Easily, the map  $f \mapsto S_f$  is a bijection from  $\{0, 1\}^\Omega$  onto  $\mathcal{P}(\Omega)$ .

Logicians often write the powerset as  $2^\Omega$ , rather than  $\{0, 1\}^\Omega$ , since all that was important about the base set  $\{0, 1\}$  was that it had 2 elements; it was not important what those elements were.  $\square$

**ENTRANCE.** Two sets  $A$  and  $B$  are **equinumerous**, or “**bijective** with each other”, if *there exists* a bijection  $A \leftrightarrow B$ . [BTWay, we use a hook-arrow to indicate an **injection**, e.g.  $A \hookrightarrow B$ , and a doublehead-arrow, e.g.  $A \leftrightarrow B$  to

indicate a **surjection**. Hence  $\leftrightarrow$  indicates a bijection.] Write the **equinumerous** relation as

$$A \asymp B.$$

Write  $A \preccurlyeq B$  if *there exists* an injection  $A \hookrightarrow B$ . Finally, let  $A \prec B$  mean that  $A \preccurlyeq B$  yet  $A \not\asymp B$ .

Easily,  $\asymp$  is an equivalence relation. [On the class of cardinalities, relation  $\preccurlyeq$  is a pre-order. Is  $\preccurlyeq$  a *partial-order*? Is  $\preccurlyeq$  a *total-order*?]

Call  $S$  **countably-infinite** or **denumerable** if  $S \asymp \mathbb{N}$ . Set  $S$  is **countable** if  $S \preccurlyeq \mathbb{N}$ , i.e.,  $S$  is bijective with some subset of  $\mathbb{N}$ . [So a countable set is either *finite* or *countably-infinite*.]  $\square$

**2a: Lemma.** Suppose sets  $P \asymp \tilde{P}$  and  $Q \asymp \tilde{Q}$ . Then  $P^Q \asymp \tilde{P}^{\tilde{Q}}$ . **Proof.** Exercise 1 soln is below.

**Proof.** By hypothesis, there are bijections  $\varepsilon: \tilde{Q} \leftrightarrow Q$  and  $\beta: \tilde{P} \leftrightarrow P$  [ $\varepsilon$  for “exponent”,  $\beta$  for “base”]. We biject  $P^Q \leftrightarrow \tilde{P}^{\tilde{Q}}$  by mapping  $f \mapsto \tilde{f}$  as in this diagram:

$$\begin{array}{ccc} Q & \xleftarrow{\varepsilon} & \tilde{Q} \\ f \downarrow & & \tilde{f} \downarrow \\ P & \xrightarrow{\beta} & \tilde{P} \end{array}$$

so  $\tilde{f} := \beta \circ f \circ \varepsilon$ . [I.e.,  $\tilde{f}(\tilde{x}) := \beta(f(\varepsilon(\tilde{x})))$  for arbitrary  $\tilde{x} \in \tilde{Q}$ .] The  $f \mapsto \tilde{f}$  mapping is (*exercise!*) a bijection. [Can you write down its inverse-map?]  $\blacklozenge$

**2b: Card-Exponentiation Lemma (CE-Lemma).** Consider any three sets  $\Omega$ ,  $B$  and  $C$ . Then  $\Omega^{B \times C} \asymp [\Omega^B]^C$ .

**Proof.** Exercise 2 soln is below.

**Proof.** Define  $\Theta: \Omega^{B \times C} \leftrightarrow [\Omega^B]^C$  by

$$\Theta(f) := \left[ c \mapsto \left[ b \mapsto f((b, c)) \right] \right].$$

Its inverse-map  $\Upsilon: [\Omega^B]^C \leftrightarrow \Omega^{B \times C}$  is

$$\Upsilon(g) := \left[ (b, c) \mapsto [g(c)](b) \right]. \quad \blacklozenge$$

**3: Countable-card Theorem.** Below,  $S$  represents an arbitrary non-void countable set.

a: An arbitrary subset of a countable set is countable. In particular, an arbitrary infinite subset of a countable set is countably-infinite.

b: Each of these is countably-infinite:  
 $\mathbb{Z}$ ,  $\mathbb{Q}$ ,  $\mathbb{N} \times \mathbb{N}$ ,  $S \times \mathbb{N}$ .

c: A union of countably many countable sets is countable.  $\diamond$

**4a: Defn.** In referring to intervals, let **LCRO** mean “Left-Closed Right-Open” and let **LORC** mean “Left-Open Right-Closed”.

Use  $\overline{\mathbb{R}} := [-\infty, +\infty] := \{-\infty\} \cup \mathbb{R} \cup \{\infty\}$  to denote the **extended reals**. In  $\overline{\mathbb{R}}$ , as an example, the set  $[-\infty, 5)$  is a LCRO-interval, and  $(7, \infty]$  is a LORC-interval. Both  $[-4, \infty]$  and  $\overline{\mathbb{R}}$  itself are closed intervals. All of these examples are *unbounded* intervals, whereas  $(-4, 7]$  is a bounded interval.

In the theorem below, the word “interval” means “non-trivial interval”. That is, we exclude 1-point closed intervals, e.g.  $[7, 7] \stackrel{\text{note}}{=} \{7\}$ , as well as the empty interval, e.g.  $(7, 7) \stackrel{\text{note}}{=} (7, 7) \stackrel{\text{note}}{=} [7, 7)$ , each of which is the emptyset. [BTWay, the emptyset is open.]  $\square$

**4b: Interval-card Theorem.** Each non-trivial sub-interval of  $\overline{\mathbb{R}}$  is equi-numerous with  $\mathbb{R}$ .  $\diamond$

**Proof.** For  $k = 1, 2$ , consider intervals  $J_k := [a_k, b_k]$ , with positive (finite) lengths  $L_k := b_k - a_k$ . Then the affine map

$$x \mapsto a_2 + \frac{L_2}{L_1} \cdot [x - a_1]$$

bijects  $J_1$  onto  $J_2$ . The same map works for two open intervals, two LCRO-intervals, or two LORC-intervals.

Let  $\mathcal{H} := \{\frac{1}{n}\}_{n=1}^{\infty}$  comprise the harmonic numbers. Define  $f: [0, 1] \rightarrow [0, 1)$  as follows.<sup>♥1</sup>

4c: For  $x \in \mathcal{H}$ : Let  $n := \frac{1}{x}$ , then map  $x \mapsto \frac{1}{n+1}$ .  
 For  $x \notin \mathcal{H}$ : Map  $x \mapsto x$ .

This  $f$  is a bijection. And  $g(x) := 1 - x$  bijects  $[0, 1)$  onto  $(0, 1]$ . Defining  $h: (0, 1] \rightarrow (0, 1)$  by rule (4c) also gives a bijection. Hence

$$[0, 1] \xrightarrow{f} [0, 1) \xrightarrow{g} (0, 1] \xrightarrow{h} (0, 1). \quad \text{Thus:}$$

All bounded intervals have the same cardinality.

**Unbounded intervals.** Extending  $\tan()$  to  $\overline{\mathbb{R}}$ , note that  $\tan$  bijects  $J := [-\frac{\pi}{2}, \frac{\pi}{2}]$  onto the closed interval  $\overline{\mathbb{R}}$ . Consequently, every (bounded or unbounded) closed/open/LORO/LCRO sub-interval of  $\overline{\mathbb{R}}$  is carried by  $\arctan$  to a corresponding subinterval of  $J$ .

<sup>♥1</sup>This (4c) uses one of our “Cantor’s–Hotel maps” from class.

Since these latter intervals are all bounded intervals, our earlier argument shows that they are all bijective with each other.♥<sup>2</sup> ♦

### 5: Cantor's diagonalization thm.

i: Firstly,  $\mathbb{N} \prec \mathbb{R}$ . Moreover, for each function  $f: \mathbb{N} \rightarrow \mathbb{R}$  there is an explicit construction of a point  $\mathcal{N}ewPt_f \in \mathbb{R}$  which is not in  $\text{Range}(f)$ .

ii: Every set  $S$  satisfies that  $S \prec \mathcal{P}(S)$ . Moreover, for each fnc  $f: S \rightarrow \mathcal{P}(S)$ , this set,

$$\mathcal{N}ewSet_f := \{z \in S \mid f(z) \not\ni z\}$$

is not in  $\text{Range}(f)$ . ♦

**Proof of (i).** Map  $x \mapsto x$  injects  $\mathbb{N}$  into  $\mathbb{R}$ , so  $\mathbb{N} \preccurlyeq \mathbb{R}$ .

Recall  $\mathbb{N} \succ \mathbb{Z}_+$  and  $\mathbb{R} \succ (0, 1)$ . Given  $g: \mathbb{Z}_+ \rightarrow (0, 1)$  we build  $\mathcal{N}ewPt_g \in (0, 1)$  with  $\mathcal{N}ewPt_g \notin \text{Range}(g)$ , as follows. For  $x \in (0, 1)$ , define digits  $\mathbf{d}_n^x$  so that

$$\sum_{n=1}^{\infty} \frac{\mathbf{d}_n^x}{10^n} = x.$$

Moreover, when  $x$  is a 10-adic rational, use the expansion which is eventually ‘9’. [Any fixed rule will work.]

If we write the decimal expansions of  $g(1), g(2), \dots$  in a two-dimensional table, then digits  $\delta_n := \mathbf{d}_n^{g(n)}$  lie along the diagonal, whence the name of the proof.

For a digit  $d$ , let

$$\bar{d} := \begin{cases} 3 & , \text{ if } d \neq 3; \\ 7 & , \text{ otherwise.} \end{cases}$$

Finally, define digit  $\alpha_n := \bar{\delta}_n$ . Then point

$$\mathcal{N}ewPt_g := \sum_{n=1}^{\infty} \frac{\alpha_n}{10^n}$$

lies in  $(0, 1)$  and [Exercise 3] is **not** in  $\text{Range}(g)$ . ♦

**Proof of (ii).** Injection  $x \mapsto \{x\}$  shows that  $S \preccurlyeq \mathcal{P}(S)$ .

To see that  $S \not\prec \mathcal{P}(S)$ , note: For each  $z \in S$ , the symmetric-difference  $\mathcal{N}ewSet_f \triangle f(z)$  owns  $z$ ; hence  $\mathcal{N}ewSet_f$  differs from  $f(z)$ . ♦

♥<sup>2</sup>BTWay, fnc  $x \mapsto e^x$  bijects  $\mathbb{R} \leftrightarrow \mathbb{R}_+$ , and  $x \mapsto e^x / [e^x + 1]$  maps  $\mathbb{R} \leftrightarrow (0, 1)$ . Indeed, each an order-preserving homeomorphism. The latter fnc is the **sigmoid function**.

## Schröder-Bernstein

Here, we examine cardinality relations between two sets,  $\mathbf{X}$  and  $\Omega$ . The below arguments do not require these sets be disjoint, but the idea is easier understand when they are. At no cost, we can arrange that  $\mathbf{X}$  be disjoint from  $\Omega$  by replacing  $\mathbf{X}$  by  $\mathbf{X} \times \{1\}$ , and replacing  $\Omega$  by  $\Omega \times \{2\}$ .

Consider a map  $g: \mathbf{X} \rightarrow \Omega$ , an injection  $h: \Omega \hookrightarrow \mathbf{X}$ , and a subset  $B \subset h(\Omega) \overset{\text{note}}{\subset} \mathbf{X}$ .

These determine a function  $\theta: \mathbf{X} \rightarrow \Omega$  by setting  $F := \mathbf{X} \setminus B$ , and defining

6.1:

$$\theta|_B := h^{-1}|_B, \text{ and } \theta|_F := g|_F.$$

Let  $\llbracket g, h: B \rrbracket$  denote this function  $\theta$ .

[So we map Forward on  $F$ , and Backward on  $B$ .]

Given  $g, h$  and  $B$  as above, with  $g$  an *injection*, say that “the set  $B$  is  **$(g, h)$ -backward-good**” if the resulting function  $\theta: \mathbf{X} \rightarrow \Omega$  of (6.1) is a *bijection*, and call  $\theta$  an “ **$(g, h)$ -good bijection**”. The corresponding forward set  $F := \mathbf{X} \setminus B$  is  **$(g, h)$ -forward-good**.

**6.2: Weak Schröder-Bernstein thm.** For sets  $\mathbf{X}$  and  $\Omega$ : If  $\mathbf{X} \preccurlyeq \Omega$  and  $\mathbf{X} \succcurlyeq \Omega$ , then  $\mathbf{X} \asymp \Omega$ . ♦

**6.3: Schröder-Bernstein Thm [S-B thm].** Fix injections

$$\mathbf{X} \xrightarrow{g} \Omega \quad \text{and} \quad \mathbf{X} \xleftarrow{h} \Omega.$$

Then there exists an  **$(g, h)$ -good bijection**.

Indeed, defining the collection of points in each set with no pre-image in the other set,

$$\widehat{X} := \mathbf{X} \setminus h(\Omega) \quad \text{and} \quad \widehat{\Omega} := \Omega \setminus g(\mathbf{X}),$$

we have this: The Smallest and Largest, in the sense of inclusion,  **$(g, h)$ -backward-good** sets are

$$6.4: \quad \mathcal{S}_{\langle g, h \rangle} := \bigcup_{n=0}^{\infty} [h \circ g]^{on} (h(\widehat{\Omega})) \quad \text{and}$$

$$6.5: \quad \mathcal{L}_{\langle g, h \rangle} := \mathbf{X} \setminus \left[ \bigcup_{n=0}^{\infty} [h \circ g]^{on} (\widehat{X}) \right],$$

resp.. In particular,  $\mathcal{S}_{\langle g, h \rangle} \subset \mathcal{L}_{\langle g, h \rangle} \subset \text{Range}(h)$ . ♦

**Proof.** Use the 4-orbit picture from class. ♦

**Bits**

Define half-open and open intervals in the reals,

$$L := [0, 1) \quad \text{and} \quad R := (0, 1] \quad \text{and} \quad J := (0, 1).$$

Recall that each dyadic rational has two binary numerals; the remaining reals each have one binary numeral. Define the all-zero and all-one names

$$\bar{0} := 000 \dots, \quad \text{and} \quad \bar{1} := 111 \dots,$$

Let EC be the set of eventually constant-1 or eventually constant-0 names.

Below, I'll use "**word**" for a *finite* string bits, e.g. "**10010**". I'll use "**name**" for a (one-sided) *infinite* string, e.g. **10111100100100100100**... (e.g, start with "**10111**" then repeat the pattern "**100**" forever.)

Let  $\text{BITS} := 2^{\mathbb{Z}^+}$  be the set of bit-sequences

$$\vec{b} = b_1 b_2 b_3 \dots, \quad \text{with each } b_j \in \{0, 1\}.$$

So  $\text{BITS} \asymp \mathcal{P}(\text{denumerable})$ .

**7a: Defn.** Define a fnc  $\text{BinOne}: (0, 1) \rightarrow \text{BITS}$  that produces the binary numeral of a point. Specifically,  $\text{BinOne}(x)$  is the unique bit-sequence  $\vec{b}$  such that

$$\sum_{n=1}^{\infty} \frac{b_n}{2^n} = x,$$

and: If  $x$  is a dyadic rational, then  $\vec{b}$  is eventually constant 1. [E.g.  $\text{BinOne}(3/4) = 10111 \dots = 10\bar{1}$ .]

Define  $\text{BinDTer}: \text{BITS} \rightarrow [0, 1]$  [Binary to Doubled-Ternary] by

$$\text{BinDTer}(\vec{b}) := \sum_{n=1}^{\infty} \frac{2b_n}{3^n}.$$

So  $\text{BinDTer}(10\bar{1})$  is the number whose base-3 numeral is **.2022**..., is  $\frac{2}{3} + \frac{1}{9} = \frac{7}{9}$ . The range of  $\text{BinDTer}$  is the famous "middle thirds" **Cantor set**.

**Exercise-4:** Both  $\text{BinOne}$  and  $\text{BinDTer}$  are injections.  $\square$

**7b: S-B corollary.** It is the case that  $\mathbb{R} \asymp \mathcal{P}(\mathbb{N})$ .  $\diamond$

**Proof.** Inject  $\mathbb{R} \hookrightarrow \text{BITS}$  by

$$\mathbb{R} \asymp (0, 1) \xrightarrow{\text{BinOne}} \text{BITS}.$$

In the other direction,

$$\text{BITS} \xrightarrow{\text{BinDTer}} [0, 1] \hookrightarrow \mathbb{R},$$

injecting  $\text{BITS} \hookrightarrow \mathbb{R}$ . Now apply Schröder-Bernstein.  $\blacklozenge$

**Explicit bijection  $\psi: \text{BITS} \hookrightarrow (0, 1]$** 

In half-open  $(0, 1]$ , let DY be this sequence of dyadics:

$$\text{DY} = \left( \frac{1}{1}, \frac{1}{2}, \frac{1}{4}, \frac{3}{4}, \frac{1}{8}, \frac{3}{8}, \frac{5}{8}, \frac{7}{8}, \frac{1}{16}, \frac{3}{16}, \frac{5}{16}, \frac{7}{16}, \dots, \frac{15}{16}, \frac{1}{32}, \frac{3}{32}, \dots, \frac{31}{32}, \dots \right).$$

Define

$$\psi(\vec{a}) := \begin{cases} \mu(\vec{a}) & , \text{ if } \vec{a} \in \text{EC}; \\ \sum_{n=1}^{\infty} \frac{a_n}{2^n} & , \text{ Otherwise} \end{cases}$$

where map  $\mu: \text{EC} \hookrightarrow \text{DY}$  is:

Eventually constant $\vec{a}$	$\mapsto$	Dyadic number $\mu(\vec{a})$
$\bar{1}$		1/1
$\bar{0}$		1/2
<b>0</b> $\bar{1}$		1/4
<b>1</b> $\bar{0}$		3/4
<b>00</b> $\bar{1}$		1/8
<b>01</b> $\bar{0}$		3/8
<b>10</b> $\bar{1}$		5/8
<b>11</b> $\bar{0}$		7/8
<b>000</b> $\bar{1}$		1/16
<b>001</b> $\bar{0}$		3/16
<b>010</b> $\bar{1}$		5/16
<b>011</b> $\bar{0}$		7/16
<b>100</b> $\bar{1}$		9/16
<b>101</b> $\bar{0}$		11/16
<b>110</b> $\bar{1}$		13/16
<b>111</b> $\bar{0}$		15/16
<b>0000</b> $\bar{1}$		1/32
<b>0001</b> $\bar{0}$		3/32
<b>0010</b> $\bar{1}$		5/32
<b>0011</b> $\bar{0}$		7/32
$\vdots$		$\vdots$
<b>1111</b> $\bar{0}$		31/32
<b>00000</b> $\bar{1}$		1/64
<b>00001</b> $\bar{1}$		3/64
<b>00010</b> $\bar{1}$		5/64
$\vdots$		$\vdots$
<b>11110</b> $\bar{0}$		61/64
<b>11111</b> $\bar{0}$		63/64
<b>000000</b> $\bar{1}$		1/128
<b>000001</b> $\bar{1}$		3/128
$\vdots$		$\vdots$

**8: Power-of-reals Thm.**

a: The plane is equi-numerous with the line.

b: For each posint  $k$ , we have that  $\mathbb{R}^k \asymp \mathbb{R}$ .

c: Cartesian-power  $\mathbb{R}^{\mathbb{N}}$  is equi-numerous with  $\mathbb{R}$ .  $\diamond$

**Pf of (a).** Map  $\text{BITS} \times \text{BITS} \xrightarrow{\mathcal{W}} \text{BITS}$  by interweaving the bits, as follows.

$$\mathcal{W}(\vec{a}, \vec{c}) := a_1 c_1 a_2 c_2 a_3 c_3 a_4 c_4 \dots$$

Pick your favorite bijection  $\mathcal{B}: \mathbb{R} \hookrightarrow \text{BITS}$ . Then

$$f(x, y) := \mathcal{B}^{-1}(\mathcal{W}(\mathcal{B}(x), \mathcal{B}(y)))$$

bijects the plane  $\mathbb{R} \times \mathbb{R}$  to the line  $\mathbb{R}$ .  $\diamond$

**Pf of (b,c).** [Note (a) is the  $k=2$  case of (b).] To prove that  $\text{BITS}^k \asymp \text{BITS}$ , we can interleave  $k$  bit-seqs.

Alternatively, producing an injection

$$\varphi: \text{BITS}^{\mathbb{N}} \hookrightarrow \text{BITS}$$

establishes (a,b,c) in one swell foop. [Trivially  $\text{BITS} \hookrightarrow \text{BITS}^{\mathbb{N}}$ , so S-B thm says  $\text{BITS}^k \asymp \text{BITS}^{\mathbb{N}} \asymp \text{BITS}$ .] We make our  $\varphi$  an actual bijection, as follows: Notice that  $\text{BITS}^{\mathbb{N}}$  can be viewed as a sequence of bit-seqs; i.e, it is a bit-quadrant, ie, a bit at each point of  $\mathbb{N} \times \mathbb{N}$ . Pick your favorite bijection showing  $\mathbb{N} \times \mathbb{N} \asymp \mathbb{N}$ , e.g, diagonal raster-scan. Then

$$\text{BITS}^{\mathbb{N}} = 2^{\mathbb{N} \times \mathbb{N}} \stackrel{\text{by (2a)}}{\asymp} 2^{\mathbb{N}} = \text{BITS}. \quad \diamond$$

**9: Continuous-fncs Thm.**

i: Firstly,  $\mathbb{R}^{\mathbb{R}} \asymp 2^{\mathbb{R}} \stackrel{\text{note}}{\asymp} \left\{ \begin{array}{l} \text{Functions only taking} \\ \text{on values 5 and 7.} \end{array} \right\}$ .

ii: Also,  $\mathbf{C}(\mathbb{R} \rightarrow \mathbb{R}) \asymp \mathbb{R} \stackrel{\text{note}}{\asymp} \{\text{Constant fncs}\}$ .  $\diamond$

Challenging: *Exercise 5.* (?)

**Pf of (i).**  $\mathbb{R}^{\mathbb{R}} \asymp [2^{\mathbb{N}}]^{\mathbb{R}} \asymp 2^{\mathbb{N} \times \mathbb{R}}$ , by CE-Lem (2b). Etc.  $\diamond$

**Pf of (ii).** An injection  $\mathbb{R} \hookrightarrow \mathbf{C}(\mathbb{R} \rightarrow \mathbb{R})$  is  $p \mapsto [x \mapsto p]$ , e.g, the number 4 maps to the constant-function-4.

Courtesy Schröder-Bernstein, then, ISTProduce an injection in the other direction. Happily,

$$\mathbf{C}(\mathbb{R} \rightarrow \mathbb{R}) \stackrel{\text{Lem (9a)}}{\asymp} \mathbb{R}^{\mathbb{Q}} \asymp [2^{\mathbb{N}}]^{\mathbb{N}} \stackrel{\text{by CE}}{\asymp} 2^{\mathbb{N} \times \mathbb{N}} \stackrel{\text{Etc.}}{\asymp} \mathbb{R}. \quad \diamond$$

**9a: Lemma.** The mapping  $\mathbf{C}(\mathbb{R} \rightarrow \mathbb{R}) \rightarrow \mathbb{R}^{\mathbb{Q}}$  defined by restriction  $f \mapsto f|_{\mathbb{Q}}$ , is an injection.  $\diamond$

**Proof.** Letting  $h := f|_{\mathbb{Q}}$ , we need to recover  $f$  from  $h$ . Given a point  $p \in \mathbb{R}$ , take a sequence  $\vec{q}$  of rationals s.t.  $q_n \rightarrow p$ . Our [unknown]  $f$  is cts at  $p$ , so

$$f(p) = \lim_{n \rightarrow \infty} f(q_n) \stackrel{\text{def}}{=} \lim_{n \rightarrow \infty} h(q_n). \quad \diamond$$

**Challenge.** Define  $\alpha: \mathbf{C}(\mathbb{R} \rightarrow \mathbb{R}) \rightarrow \mathbb{R}^{\mathbb{Q}}$  by  $\alpha(f) := f|_{\mathbb{Q}}$ . Is  $\alpha()$  surjective?

For  $f \in \mathbf{C}(\mathbb{R} \rightarrow \mathbb{R})$ , restriction  $f|_{\mathbb{Q}}$  lies in  $\mathbb{R}^{\mathbb{Q}}$ ; but is it necessarily a continuous map  $\mathbb{Q} \rightarrow \mathbb{R}$ ?

If you think “yes”, then: Does the above  $\alpha()$  map onto  $\mathbf{C}(\mathbb{Q} \rightarrow \mathbb{R})$ ?  $\square$

## Algebraic numbers

A complex number  $\gamma$  is **algebraic** if it is a root of some non-*zip* intpoly [equiv., ratpoly]  $f$ . Thus

$$\alpha := \sqrt[5]{19} \quad \text{and} \quad \beta := [1 - \sqrt{13}]/6$$

are algebraic numbers, since  $\alpha$  is a root of  $x^5 - 19$ , and  $\beta$  is a root of  $3x^2 - x - 1$ . Evidently each rational number  $P/Q$  is algebraic, since it is a root of intpoly  $Qx - P$ .

Each algebraic number  $\gamma$  has an associated posint called its **degree**, written  $\text{Deg}(\gamma)$ . Writing  $\mathbf{d} := \text{Deg}(\gamma)$ , then  $\gamma$  is a root of some degree- $\mathbf{d}$  intpoly, but is the root of *no lower-degree* [non-*zip*] intpoly.

The rationals are precisely those numbers of degree 1. The above  $\alpha$  has  $\text{Deg}(\alpha) \leq 5$ . The above  $\beta$  has  $\text{Deg}(\beta) = 2$ , since  $\sqrt{13}$  is irrational.

Use  $\mathbb{A}$  for the set of algebraic numbers in  $\mathbb{C}$ . We see that  $\mathbb{A}$  is stratified into a *hierarchy* by degree. The numbers in the complement,  $\mathbb{C} \setminus \mathbb{A}$ , *transcend* this hierarchy so –not surprisingly– each such number is said to be **transcendental**. Although this is not obvious, each of these three numbers

$$\pi, \quad e, \quad \tau := \sum_{n=1}^{\infty} \frac{1}{b_n}, \quad \text{where } b_n := 2^{n!},$$

is transcendental.<sup>♥3</sup>

We define the **degree** of a transcendental number to be  $\infty$ . That is to say, the degree of a number  $\gamma \in \mathbb{C}$  is the *infimum* of numbers  $d \in [1.. \infty)$  such that  $\gamma$  is a zero of some degree- $d$  intpoly.

*Defn.* Algebraists use notation  $\mathbb{Z}[x]$  for the set of  $\mathbb{Z}$ -coefficient polynomials written using variable “ $x$ ”. They use  $\mathbb{Q}[x]$  for the set of rational-coeff polynomials.  $\square$

**11: Lemma.** *Sets  $\mathbb{Z}[x]$  and  $\mathbb{Q}[x]$  are equi-numerous, and each is countably-infinite.*  $\diamond$

*Pf.* Certainly  $\mathbb{Q}[x] \asymp \mathbb{Z}[x] \asymp \mathbb{N}[x]$ , since  $\mathbb{Q} \asymp \mathbb{Z} \asymp \mathbb{N}$ .

Let  $p_k$  be the  $k^{\text{th}}$ -prime; so  $p_0=2$ ,  $p_1=3$ ,  $p_2=5$ ,  $\dots$ . A  $\mathbb{N}$ -coefficient polynomial can be written uniquely as

$$\dagger: \quad c_0 + c_1x + c_2x^2 + c_3x^3 + \dots,$$

where  $\infty$ -sequence  $\vec{c}$  is eventually-constant-zero. Mapping  $(\dagger)$  to

$$\ddagger: \quad \prod_{k=0}^{\infty} [p_k]^{c_k},$$

is well-defined [since  $\vec{c}$  is eventually-const-zero] and is an bijection  $\mathbb{N}[x] \hookrightarrow \mathbb{Z}_+$ . [E.g.,  $\text{Zip}$  maps to  $2^0 \cdot 3^0 \cdot 5^0 \dots = 1$ . And  $3x^2 + x^4$  maps to  $p_2^3 \cdot p_4^1 = 5^3 \cdot 11 = 1375$ .]  $\blacklozenge$

**12: Algebraic-numbers Thm.** *The set  $\mathbb{A}$  of algebraic numbers is denumerable.*  $\diamond$

*Pf.* Each [non-*zip*] intpoly has only finitely-many roots. And Lemma (11) asserts only countably many intpolys. Thus  $\mathbb{A}$  is a countable union of countable sets, hence is countable, courtesy (3), the Countable-card theorem.  $\blacklozenge$

(And now, the Appendices! See next page.)

<sup>♥3</sup>Such a  $\tau$  is called a *Liouville number*. There is an explanation of Liouville numbers on my Teaching Page.

## §A Schröder-Bernstein computation

*S-B Challenge.* Let's use **Schröder-Bernstein** to construct a bijection  $\theta: \mathbb{Z}_+ \rightarrow \mathbb{Q}_+$ . [Since  $\mathbb{Z}_+$  and  $\mathbb{Q}_+$  are not disjoint, I'll use **blue** for  $\mathbb{Z}_+$  and its elements, and use **reddish** colors for  $\mathbb{Q}_+$  and its elements.] Define the Divides map  $D: \mathbb{Z}_+ \rightarrow \mathbb{Q}_+$  by

$$D(n) := n/3.$$

Easily,  $D$  is well-defined, and is an injection.

Each positive rational can be uniquely written  $p/q$ , where  $p \perp q$  are posints. Define  $R: \mathbb{Q}_+ \rightarrow \mathbb{Z}_+$  [it applies to Ratios] by

$$R(p/q) := 2^{p-1} \cdot 3^{q-1}.$$

[For example,  $R(1/1) = 2^{1-1} \cdot 3^{1-1} = 2^0 3^0 = 1$ . And  $R(7/3) = 2^{7-1} \cdot 3^{3-1} = 576$ .] This  $R$  is well-defined, and is injective because prime-factorization is *unique*.

Let  $\theta: \mathbb{Z}_+ \rightarrow \mathbb{Q}_+$  be the  $(D, R)$ -good bijection with the *smallest* backward set. [So  $\theta$  uses  $R^{-1}$  *only* on those  $(D, R)$ -orbits that start in  $\mathbb{Q}_+$ .]

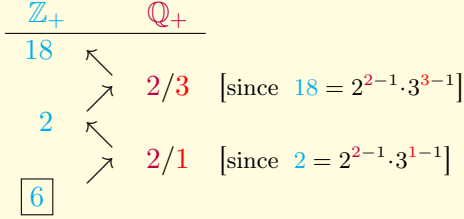
Compute  $\theta(18)$ .      Compute  $\theta(4)$ . □

*Want to think about it first? Good idea!*

*WARNING: A soln is on the next page.*

**Looking into the Past.** When does an  $n \in \mathbb{Z}_+$  have an  $R$ -preimage? Exactly when  $n$  factors as  $2^{p-1} \cdot 3^{q-1}$  with posints  $p$  and  $q$  coprime to each other.

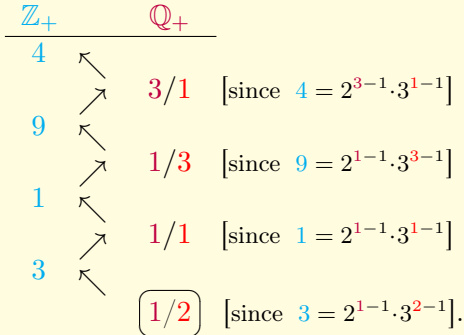
**Backtracing 18.** This gives



This 6 has no  $R$ -preimage. For although 6 is a product of powers of 2 and 3, the exponents are *not* coprime:  $6 = 2^{2-1} \cdot 3^{2-1}$ , and 2 is not coprime to 2. Hence:

$$\dagger: \quad \theta(18) \stackrel{\text{Fwd}}{=} D(18) = 18/3 = 6/1.$$

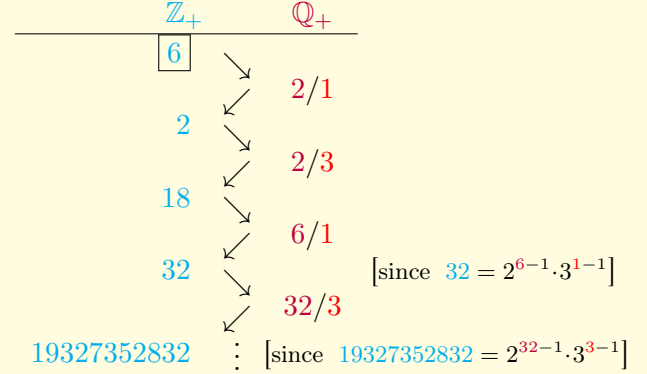
**Backtracing 4.** We compute:



Backtracing terminates with  $\frac{1}{2}$ , since  $3 \cdot \frac{1}{2}$  is not in  $\mathbb{Z}_+$  i.e.,  $\frac{1}{2} \notin \text{Range}(D)$ . Thus

$$\dagger: \quad \theta(4) \stackrel{\text{Bwd}}{=} R^{-1}(4) = 3/1 = 3. \quad \blacklozenge$$

**Forward tracing 6.** As an example, let's follow the  $(D, R)$ -orbit of 6 a bit further:



**Discussion.** Whether the S-B algorithm is constructive depends one's definition of "constructive". How to constructively tell if an orbit has no beginning?



*S-B exer.* Let  $\mathbf{X} := [3.. \infty)$  and  $\mathbf{\Omega} := [4.. \infty)$ . Define

$$g: \mathbf{X} \rightarrow \mathbf{\Omega} \quad \text{by} \quad g(n) := n + 7 \quad \text{and} \\ h: \mathbf{\Omega} \rightarrow \mathbf{X} \quad \text{by} \quad h(k) := 3k + [-1]^k.$$

With  $\theta$  the (unique, in this case) S-B map  $\mathbf{X} \leftrightarrow \mathbf{\Omega}$ , compute:  $\theta(10, 11, \dots, 19)$  and  $\theta(50, 51, \dots, 59)$ .

Compute  $\theta^{-1}()$  of various numbers.

What is the smallest  $k \in \mathbf{\Omega}$  whose backwards  $g$ - $h$ -orbit has length 5?  $\square$

*Two computations.* To help out:

Back-tracing the pair-orbit:

X	Omega
206	
<---	69
62 --->	
<---	21
14 --->	
<---	5

Thus...

$$\theta(206) = H^{-1}(206) = 69.$$

=====

Back-tracing the pair-orbit:

X	Omega
368	
<---	123
116 --->	
<---	39
32 --->	
<---	11
4 --->	

Consequently...

$$\theta(368) = G(368) = 375.$$

*$\theta$  backwards.* Back-tracing from the  $\mathbf{\Omega}$  side:

Omega	X
375	
<---	368
123 --->	
<---	116
39 --->	
<---	32
11 --->	
<---	4

Hence...

$$\theta^{-1}(375) = G^{-1}(375) = 368.$$

---

Again from  $\mathbf{\Omega}$ , back-tracing:

Omega	X
105	
<---	98
33 --->	
<---	26
9 --->	

We ended on the  $\mathbf{\Omega}$  side, so

$$\theta^{-1}(105) = H(105) = 314.$$

Is  $\pi$  is on the way?  $\blacklozenge$

## §B Appendix: AC $\implies$ WO

### Orders

Below, **order** means *strict total-order*. For an order-symbol  $<, \prec, \sqsubset$ , use  $\leq, \preceq, \sqsubseteq$ , for the non-strict versions.

Fixing an order  $<$  on set  $\mathbf{X}$  and a  $p \in \mathbf{X}$ , let

$$\mathbf{X}^{< p} := \{x \in \mathbf{X} \mid x < p\},$$

and analogously for  $\mathbf{X}^{\leq p}$ . Subset  $\mathbf{I} \subset \mathbf{X}$  is an “**initial segment** of  $\mathbf{X}$ ” if  $[\forall s \in \mathbf{I}, \forall x \in \mathbf{X}: x < s \implies x \in \mathbf{I}]$ . Every non-void init-seg is of form  $\mathbf{X}^{< p}$  or  $\mathbf{X}^{\leq p}$ .

For orders  $\langle \mathbf{X}, < \rangle$  and  $\langle \mathbf{\Omega}, \prec \rangle$ , an “**order-embedding**,  $f$ , of  $\mathbf{X}$  into  $\mathbf{\Omega}$ ”, written  $f: \mathbf{X} \xrightarrow{\text{emb}} \mathbf{\Omega}$ , means

$$\forall a, b \in \mathbf{X}: a < b \text{ IFF } \varphi(a) \prec \varphi(b).$$

[Another name is an “**into-isomorphism**”.]

Write  $f: \mathbf{X} \xrightarrow{\text{init}} \mathbf{\Omega}$  if  $f: \mathbf{X} \xrightarrow{\text{emb}} \mathbf{\Omega}$  and  $\text{Range}(f)$  is an *initial-segment* of  $\mathbf{\Omega}$ .

**14a: Prop'n.** For well-orders  $\langle \mathbf{X}, < \rangle$  and  $\langle \mathbf{\Omega}, \prec \rangle$ , suppose  $\varphi: \mathbf{X} \xrightarrow{\text{init}} \mathbf{\Omega}$  and  $\lambda: \mathbf{X} \xrightarrow{\text{init}} \mathbf{\Omega}$ . Then  $\varphi = \lambda$ .  $\diamond$

**Pf.** Assuming  $\varphi \neq \lambda$ , let  $\mathbf{t} \in \mathbf{X}$  be the *smallest*  $\mathbf{X}$ -value s.t, WLOG,  $\lambda(\mathbf{t}) \succ \varphi(\mathbf{t}) =: \tau$ . For each  $x < \mathbf{t}$ , then,  $\lambda(x) = \varphi(x) \prec \tau \prec \lambda(\mathbf{t})$ . Thus  $\lambda()$  skips over  $\tau$ , hence is not an init-seg map.  $\otimes$

**14b: Lemma.** Fix well-orders  $\langle \mathbf{X}, < \rangle$  and  $\langle \mathbf{\Omega}, \prec \rangle$ . If they are order-isomorphic, then the isomorphism is unique.

If not, then exactly one of them is ord-iso to a subset the other. Moreover, it admits an ord-iso to an *initial-segment*, and this  $\xrightarrow{\text{init}}$  map is unique.  $\diamond$

**Pf.** Let  $\mathcal{C}$  be the set of  $p \in \mathbf{X}$  for which init-seg  $\mathbf{X}^{\leq p}$  admits a map  $f_p: \mathbf{X}^{\leq p} \xrightarrow{\text{init}} \mathbf{\Omega}$ . For  $s > p$ , both in  $\mathcal{C}$ , our Prop'n (14a) implies that the restriction of  $f_s$  to  $\mathbf{X}^{\leq p}$  equals  $f_p$ . Consequently, the union

$$\varphi := \bigcup_{p \in \mathcal{C}} f_p$$

is a well-defined map into  $\mathbf{\Omega}$ . Its domain is initial-segment

$$\mathbf{I} := \bigcup_{p \in \mathcal{C}} \mathbf{X}^{\leq p}.$$

This  $\varphi$  is an ord-iso, since each  $f_p$  is, and maps onto  $\mathbf{\Omega}$ -init-seg

$$\mathbf{\Lambda} := \bigcup_{p \in \mathcal{C}} \text{Range}(f_p).$$

**Which direction?** If  $\mathbf{I}$  equals  $\mathbf{X}$ , then  $\varphi: \mathbf{X} \xrightarrow{\text{init}} \mathbf{\Omega}$ .

Otherwise, let  $\mathbf{s} := \text{Min}(\mathbf{X} \setminus \mathbf{I})$ . Could  $\mathbf{\Lambda}$  fail to be all of  $\mathbf{\Omega}$ ? *No!*, since otherwise we could extend  $\varphi$  by mapping  $\mathbf{s}$  to  $\text{Min}(\mathbf{\Omega} \setminus \mathbf{\Lambda})$ . Hence  $\varphi^{-1}: \mathbf{\Omega} \xrightarrow{\text{init}} \mathbf{X}$ .  $\blacklozenge$

**14c: Corollary.** On the proper-class of Well-Order-types, relation  $\xrightarrow{\text{init}}$  is a [lax, i.e, non-strict] well-order.  $\diamond$

**Well-ordering Axiom.** The WO Axiom states that each set admits a well-order.  $\square$

**15: WOA  $\implies$  AC thm.** Assuming WO Axiom, each collection  $\mathcal{C}$  of non-void sets, admits a choice fnc.  $\diamond$

**Proof.** Let  $<$  be a well-order on  $\mathbf{U} := \bigcup(\mathcal{C})$ . This engenders choice-fnc  $A \mapsto \text{Min}^<(A)$ , for each  $A \in \mathcal{C}$ .  $\blacklozenge$

### Zermelo's pf AC $\Rightarrow$ Well-ordering Principle

In 1904, Ernst Zermelo proved the then-surprising result that AC implies WOAxiom.

**16a: Shy-function.** On a set  $\mathbf{X}$ , let  $\mathfrak{M} := 2^{\mathbf{X}} \setminus \{\mathbf{X}\}$  be the collection of *proper* subsets. The Axiom-of-Choice gives the existence of a *shy-fnc*  $\mathcal{Y}:\mathfrak{M} \rightarrow \mathbf{X}$  satisfying

$$\dagger: \quad \forall S \in \mathfrak{M}: \quad \mathcal{Y}(S) \in \mathbf{X} \setminus S,$$

[The shy-fnc picks an  $\mathbf{X}$ -element  $\mathcal{Y}(S)$  that *avoids*  $S$ .] A shy-fnc comes from AC applied to collection  $\{\mathbf{X} \setminus S\}_{S \in \mathfrak{M}}$  of non-void sets.

Henceforth, there is a fixed a shy-fnc  $\mathcal{Y}$  on  $\mathbf{X}$ .

On a subset  $S \subset \mathbf{X}$ , a well-order  $\prec$  is “*good* on  $S$ ” [or “*pair*  $\langle S, \prec \rangle$  is good”] if

$$\dagger: \quad \forall t \in S: \quad \mathcal{Y}(S^{\prec t}) = t. \quad \square$$

**16b: Obs.** Fix a good  $\langle S, \prec \rangle$ . For each proper  $\prec$ -init-seg  $I \subsetneq S$ , let  $\mathbf{t} := \text{Min}^{\prec}(S \setminus I)$ . Thus  $I = S^{\prec \mathbf{t}}$ . Hence

$$\pounds: \quad \text{Min}^{\prec}(S \setminus I) = \mathcal{Y}(I). \quad \square$$

**16c: Shy lemma.** For subsets  $S, T \subset \mathbf{X}$ , suppose pairs  $\langle S, \prec \rangle$  and  $\langle T, \prec \rangle$  are each good. Then either  $S \subset T$  or  $T \subset S$ .

When  $S \subset T$ , then  $S$  is a  $\prec$ -initial-segment. Further,  $\prec$  equals  $\prec|_S$ ; the  $\prec$ -order restricted to  $S$ .

[IOWords,  $\langle S, \prec \rangle \xrightarrow{\text{init}} \langle T, \prec \rangle$  via the identity-map.]  $\diamond$

**16d: Prelim.** A subset  $J \subset \mathbf{X}$  is *mutual* if  $J \subset S \cap T$ , together with

$J$  is init-seg w.r.t  $\prec$  and w.r.t  $\prec$ , and orders  $\prec$  and  $\prec$  agree on  $J$ .  $\square$

**Pf.** Let  $\mathcal{C}$  comprise those  $p \in S$  s.t.  $S^{\prec p}$  is *mutual*. Automatically, the union  $I := \bigcup_{p \in \mathcal{C}} S^{\prec p}$  is mutual. [We don't need this, but note  $\mathcal{C} = I$ .]

**Inclusion.** If  $I \subsetneq S \cap T$ , then  $(\pounds)$  gives

$$\text{Min}^{\prec}(S \setminus I) = \mathcal{Y}(I) = \text{Min}^{\prec}(T \setminus I) \stackrel{\text{note}}{\in} S \cap T.$$

With  $\mathbf{y} := \mathcal{Y}(I)$ , then  $S^{\prec \mathbf{y}} = I \sqcup \{\mathbf{y}\} = T^{\prec \mathbf{y}}$ . Orders  $\prec$  and  $\prec$  agree on  $I \sqcup \{\mathbf{y}\}$ , yielding  $\otimes$  that  $\mathbf{y}$  is in  $I \stackrel{\text{recall}}{=} \bigcup_{p \in \mathcal{C}} S^{\prec p}$ .

If  $S = I$ , then  $S \subset T$  is a  $\prec$ -init-segment on which orders  $\prec$  and  $\prec$  agree. And if  $T = I$ , then  $T \subset S$  is a  $\prec$ -init-segment on which orders  $\prec$  and  $\prec$  agree.  $\blacklozenge$

**Prelim and Caveat.** Consider  $\mathcal{C}$ , a collection of  $\langle S, \prec_S \rangle$  pairs with  $\prec_S$  a partial-order on  $S \subset \mathbf{X}$ . This  $\mathcal{C}$  is *consistent* if for each  $\langle S, \prec_S \rangle$  and  $\langle T, \prec_T \rangle$ , partial-orders  $\prec_S$  and  $\prec_T$  agree on  $S \cap T$ . When, further, always either  $S \subset T$  or  $T \subset S$ , then  $\mathcal{C}$  is *nested*.

Define relation

$$*: \quad \prec := \bigcup_{\langle S, \prec_S \rangle \in \mathcal{C}} \prec_S \quad \text{on set} \quad \mathbf{U} := \bigcup_{\langle S, \prec_S \rangle \in \mathcal{C}} S.$$

When  $\mathcal{C}$  *consistent*, then  $\prec$  is a partial-order [exercise]. If each  $\prec_S$  is a *total*-order, and  $\mathcal{C}$  is *nested*, then  $\prec$  is a *total*-order [exercise].

If, in addition, each  $\prec_S$  is a well-order, must  $\prec$  be a WOrder? *No!* Let  $S_n := [-n .. \infty) \subset \mathbb{Z}$ , for  $n = 1, 2, \dots$ , with order  $\prec_n$  being  $\prec|_{S_n}$ . The  $(*)$ -union gives relation  $\prec$  on  $\mathbf{U} = \mathbb{Z}$ ; not a well-order.  $\square$

**17: Zermelo's W-O Thm.** If a set  $\mathbf{X}$  admits a shy-fnc, then  $\mathbf{X}$  admits a well-order.  $\diamond$

**Pf.** Let  $\mathcal{C}$  comprise all good pairs  $\langle S, \prec_S \rangle$ , where  $S \subset \mathbf{X}$ , and use  $(*)$  to define relation  $\prec$  on set  $\mathbf{U}$ . Our  $\mathcal{C}$  is nested, courtesy the Shy lemma; hence  $\prec$  is a total-order.

**$\prec$  is a well-order.** Fix a non-void subset  $B \subset \mathbf{U}$ .

For  $j=1,2$ , consider pairs  $\langle S_j, \prec_j \rangle$  having intersection  $B \cap S_j$  non-void. Let  $\mathbf{s}_j$  be the  $\prec_j$ -min of  $B \cap S_j$ . By the Shy lemma, WLOG  $S_1$  is a  $\prec_2$ -init-seg of  $S_2$ ; thus  $\mathbf{s}_2 = \mathbf{s}_1$ . Hence  $\mathbf{s}_1$  is  $\text{Min}^{\prec}(B \cap \mathbf{U})$ .

**Well-order  $\langle \mathbf{U}, \prec \rangle$  is good.** Fix a  $t \in \mathbf{U}$ . There exists a good  $\langle S, \prec \rangle$  with  $t \in S$ . The Shy lemma implies  $S$  is a  $\prec$ -init-seg. Thus  $\mathbf{U}^{\prec t} = S^{\prec t} = t$ .

**U is everything.** If  $\mathbf{U} \subsetneq \mathbf{X}$ , let  $\mathbf{y} := \mathcal{Y}(\mathbf{U})$ . Extend  $<$  to well-order  $\hat{<}$  on  $\hat{\mathbf{U}} := \mathbf{U} \sqcup \{\mathbf{y}\}$  by defining  $u \hat{<} \mathbf{y}$  for each  $u \in \mathbf{U}$ . Easily,  $\langle \hat{\mathbf{U}}, \hat{<} \rangle$  is good, contradicting that  $\mathcal{C}$  comprised *all* good pairs.  $\blacklozenge$

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