

## 2003 A6 Putnam Prob and Soln

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### 2003 Putnam

**1: A6.2003.** For an  $S \subset \mathbb{N}$ , let  $\mathbf{r}_S(n)$  denote the number of ordered pairs  $(s, t)$  with  $\{s, t\} \subset S$ ,  $s \neq t$ , and  $s+t = n$ . [This  $(s, t)$  is an *S-representation* of  $n$ .] Is it possible to partition  $\mathbb{N} = A \sqcup B$  in such a way that  $\mathbf{r}_A(n) = \mathbf{r}_B(n)$  for all natnums  $n$ ?  $\diamond$

*Soln A6.2003 (JK).* Yes; the *Morse minimal sequence* is the unique solution.

Rename the sets from  $A, B$  to  $E, D$  (for Even and odd), and require that  $E \ni 0$ .

**Uniqueness.** Fix a soln  $\mathbb{N} = E \sqcup D$ . For each posint  $n$ , we will show that  $E' := E \cap [0..n]$  and  $D' := D \cap [0..n]$  determine where  $n$  goes. Recall  $D \not\ni 0$ , so  $\mathbf{r}_D(n) = \mathbf{r}_{D'}(n)$ . Thus

$$[D \ni n] \iff [\mathbf{r}_{E'}(n) = \mathbf{r}_{D'}(n)].$$

Indeed, inequality in (\*), above, necessarily means that  $\mathbf{r}_{D'}(n)$  equals  $2 + \mathbf{r}_{E'}(n)$ . So we give  $E$  two more representations,  $(0, n)$  and  $(n, 0)$ , by placing  $n$  in  $E$ .

**Construction.** For  $n \in \mathbb{N}$ , let  $\#n$  denote the number of 1-bits in  $\text{BinaryNumeral}(n)$ . E.g.  $\#(23) = 4$ , since  $23 = 10111$ . Define

$$\begin{aligned} 1a: \quad E &:= \{n \in \mathbb{N} \mid \#n \text{ is even}\} \\ &D := \{n \in \mathbb{N} \mid \#n \text{ is odd}\}. \end{aligned}$$

**Involution.** Each subset  $S \subset \mathbb{N}$  engenders

$$\widehat{S} := \{(s, t) \mid s, t \in S \text{ with } s \neq t\}.$$

Given natnums  $s \neq t$ , write them in binary as

$$\begin{aligned} 1b: \quad s &= b_K b_{K-1} \dots b_4 b_3 b_2 b_1 b_0 \quad \text{and} \\ &t = c_K c_{K-1} \dots c_4 c_3 c_2 c_1 c_0. \end{aligned}$$

(Possibly  $b_K$  or  $c_K$  is 0.) Let  $\ell$  be the smallest index st.  $b_\ell \neq c_\ell$ . Now *exchange* these two bits, defining a new pair  $(s', t')$ ; for example, if  $\ell = 2$  then

$$\begin{aligned} 1c: \quad s' &= b_K b_{K-1} \dots b_4 b_3 c_2 b_1 b_0 \quad \text{and} \\ &t' = c_K c_{K-1} \dots c_4 c_3 b_2 c_1 c_0. \end{aligned}$$

The map  $(s, t) \xrightarrow{f} (s', t')$  is an involution on  $\widehat{\mathbb{N}}$ . Moreover, it bijects  $\widehat{E}$  with  $\widehat{D}$ . Lastly,  $f$  is sum-preserving;  $s' + t' = s + t$ . Thus  $\mathbf{r}_E(n) = \mathbf{r}_D(n)$ .  $\diamond$

*Post provem.* Define a 1-sided-infinite bit-sequence  $b_0 b_1 b_2 \dots$ , where  $b_n$  is the parity of  $\#n$ . This gives

$$0110\ 1001\ 1001\ 0110\ 1\dots,$$

the *Morse minimal sequence*.  $\square$

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