

## Euclidean algorithm in Lightning-Bolt form

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The Euclidean algorithm, **EU**, is often presented by a series of equations. I have found the following table-form convenient, both because it organises the computation, and gives a name to each number in the table. Because the update-rule follows the shape of a lightning-bolt, I call it the LBolt algorithm.

Henceforth, all variables are *integers* unless explicitly stated otherwise. Given integers  $r_0$  and  $r_1$  (for the time being, assume each is positive) we will compute  $\mathcal{G} := \text{GCD}(r_0, r_1)$  as well as a pair  $S, T$  of **Bézout multipliers** satisfying

$$1: \quad \mathcal{G} = S \cdot r_0 + T \cdot r_1.$$

[There is a one-parameter family of Bézout-pairs; the algorithm will compute a particular pair.] I'll explain via an example. Suppose we want the GCD of  $r_0 := 114$  and  $r_1 := 33$ . Then initialize the table as:

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	114	—	1	0
1	33		0	1

In order to compute a **BP** (Bézout-Pair), we'll need

$$1': \quad r_n = s_n \cdot 114 + t_n \cdot 33$$

to hold, for *every*  $n$ . Notice that it *already holds*, trivially, for  $n=0$  and  $n=1$ .

At stage  $n$ , divide  $r_n$  into  $r_{n-1}$  to get a quotient,  $q_n$ , and a remainder,  $r_{n+1}$ . That is,

$$2: \quad r_{n-1} = [q_n r_n] + r_{n+1}.$$

Now use this value of  $q_n$  to update three columns:

$$3: \quad \begin{aligned} r_{n+1} &= r_{n-1} - q_n r_n; \\ s_{n+1} &:= s_{n-1} - q_n s_n; \\ t_{n+1} &:= t_{n-1} - q_n t_n. \end{aligned}$$

Doing this for  $n=1$  gives

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	114	—	1	0
1	33	3	0	1
2	15		1	-3

Continue until you get a “0” in the  $r$ -column; I'll compute the resulting “quotient” and write “ $\infty$ ” in the  $q$ -column, obtaining

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	114	—	1	0
1	33	3	0	1
2	15	2	1	-3
3	3	5	-2	7
4	0	$\infty$	11	-38

The **GCD-row** [shown here red and italicized] is the row *above* the “ $\infty$ -row”. The numbers we sought lie in the GCD-row. In this instance,  $\mathcal{G} = r_3$ ,  $S = s_3$  and  $T = t_3$ . And indeed,

$$3 = -2 \cdot 114 + 7 \cdot 33.$$

**Why the extra row?** You wonder “Why bother to compute  $s_4$  and  $t_4$ ?” It isn't necessary, but they provide verification-data. Consider finding  $(\mathcal{G}, S, T)$  when  $r_0 := 98$  and  $r_1 := 51$ . Initialize:

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	98	—	1	0
1	51		0	1

Now compute...

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	98	—	1	0
1	51	1	0	1
2	47	1	1	-1
3	4	11	-1	2
4	3	1	12	-23
5	1	3	-13	25
6	0	$\infty$	51	-98

This  $r_5$ , which is **1**, is indeed  $\text{GCD}(98, 51)$ . And

$$1 = [-13] \cdot 98 + 25 \cdot 51.$$

Now examine the  $\infty$ -row; here, the 6<sup>th</sup> row. Note that  $s_6$  equals  $r_1$  upto  $\pm$ . And  $t_6$  equals  $r_0$  upto  $\pm$ .

In general, letting  $\mathcal{G} := \text{GCD}(r_0, r_1)$ , this “extra” row satisfies<sup>♥1</sup> that

$$4: \quad s_{N+1} \cdot \mathcal{G} = r_1 \cdot [-1]^{N+1} \quad \text{and} \quad t_{N+1} \cdot \mathcal{G} = r_0 \cdot [-1]^N.$$

<sup>♥1</sup>This is stated formally, and proven, in (9c), further below.

If you made a computational error earlier in the table, a glance at this  $[N+1]^{th}$ -row will usually shout “Error!”.

**Convention.** Depending on context, agree to use “GCD-row” to mean both its index, and its contents. E.g, for the preceding LBolt table, expression “Let  $N := \text{GCD-row}$ ” makes  $N = 5$ . I might also say “In the GCD-row, the  $t$ -value is 25.”

**Related pamphlets.** Our *Teaching page*

<http://www.math.ufl.edu/~squash/teaching.html>

has link “*practice sheet for the LBolt alg*” with pre-made tables.

There, too, is link “*Algorithms in Number Theory*” which uses LBolt iteratively to compute the  $\mathcal{G} := \text{GCD}(M_1, M_2, \dots, M_L)$  of a *list* of integers, computing also a Bézout multipliers  $S_1, S_2, \dots, S_L$  st.

$$5: \quad \sum_{\ell=1}^L S_{\ell} M_{\ell} = \mathcal{G}.$$

We call  $\vec{S} := (S_1, \dots, S_L)$  a *Bézout tuple* for the given tuple  $\vec{M} := (M_1, \dots, M_L)$ .

*Exer: Fix an  $L$ -tuple  $\vec{M}$  which is not the all-zero tuple. Prove that the set of Bézout tuples for  $\vec{M}$  is  $[L-1]$ -dimensional.*

The 2<sup>nd</sup> page of “*Algorithms in NT*” describes an algorithm for solving linear congruences such as  $33x \equiv_{114} 18$ , and has a worked-example.

## Proving the Euclidean Alg. works

I’ll leave this as an **Exer: The Euclidean-Alg always halts.**

Define the divisor and common-divisor sets,

$$\mathcal{D}(K) := \{d \in \mathbb{Z} \mid d \mid K\} \quad \text{and} \\ \mathcal{C}(K, N) := \mathcal{D}(K) \cap \mathcal{D}(N).$$

[Below, “LC” stands for “Linear Combination”].

**6: LC Lemma.** Consider integers  $\alpha, \beta, \gamma, M$  such that

$$6a: \quad \alpha + [M \cdot \beta] = \gamma.$$

Then

$$*: \quad \mathcal{C}(\alpha, \beta) = \mathcal{C}(\beta, \gamma). \quad \diamond$$

**Proof.** Each  $d \in \mathcal{C}(\alpha, \beta)$  necessarily divides  $\alpha + [M\beta]$ , since  $M \in \mathbb{Z}$ . Thus  $\mathcal{C}(\alpha, \beta) \subset \mathcal{D}(\gamma)$ . By its definition,  $\mathcal{C}(\alpha, \beta) \subset \mathcal{D}(\beta)$ . Consequently

$$6b: \quad \mathcal{C}(\alpha, \beta) \subset \mathcal{C}(\beta, \gamma).$$

OTOHand, we can rewrite (6a) as

$$\gamma + [-M \cdot \beta] = \alpha.$$

The above reasoning hands us

$$6c: \quad \mathcal{C}(\alpha, \beta) \supset \mathcal{C}(\beta, \gamma).$$

This, together with (6b), yields (\*). ♦

**6d: Corollary.** Consider an LBolt seeded with integers  $r_0$  and  $r_1$ . Then  $\mathcal{C}(r_k, r_{k+1}) = \mathcal{C}(r_0, r_1)$ , for each index  $k$ . Consequently,

$$6e: \quad \text{GCD}(r_k, r_{k+1}) = \text{GCD}(r_0, r_1).$$

Letting  $N$  be the GCD-row index, then,

$$6f: \quad r_N = \text{GCD}(r_0, r_1),$$

since  $r_{N+1}$  is zero. ♦

**7: Bézout Lemma.** Consider an *LBolt* seeded with integers  $r_0$  and  $r_1$ . For each  $k$ , then,

$$\mathbf{B}(k): \quad r_k = [s_k r_0] + [t_k r_1]$$

holds. I'll refer to assertion  $[\forall k \in \mathbb{N}: \mathbf{B}(k)]$  as the **Bézout row-property** or **LBolt row-property**.

With  $N := \text{GCD-index}$ , consequently,

$$7a: \quad \text{GCD}(r_0, r_1) = [s_N r_0] + [t_N r_1]. \quad \diamond$$

**Proof.** The *LBolt*-seeding gives  $\mathbf{B}(0)$  and  $\mathbf{B}(1)$ .

Now fix a posint  $n$  st.  $\mathbf{B}(n-1)$  and  $\mathbf{B}(n)$ . Courtesy update rule (3),

$$\begin{aligned} & s_{n+1}r_0 + t_{n+1}r_1 \\ &= [s_{n-1} - q_n s_n] \cdot r_0 + [t_{n-1} - q_n t_n] \cdot r_1 \\ &= [s_{n-1}r_0 + t_{n-1}r_1] - q_n \cdot [s_n r_0 + t_n r_1], \end{aligned}$$

since multiplication distributes-over addition. Assertions  $\mathbf{B}(n-1)$  and  $\mathbf{B}(n)$  now give us that

$$s_{n+1}r_0 + t_{n+1}r_1 = r_{n-1} - q_n \cdot r_n$$

which, by update (3), equals  $r_{n+1}$ . We've thus inductively established

$$\forall k \geq 1: \quad [\mathbf{B}(k-1) \ \& \ \mathbf{B}(k)] \implies \mathbf{B}(k+1). \quad \blacklozenge$$

### Alternate initialization

Consider an LBolt seeded with integers  $r_0$  and  $r_1$ . Define matrices

$$M_n := \begin{bmatrix} s_n & t_n \\ s_{n+1} & t_{n+1} \end{bmatrix} \quad \text{and} \quad R_n := \begin{bmatrix} r_n \\ r_{n+1} \end{bmatrix}.$$

Up till now, our initialization matrix  $M_0$  has the identity matrix  $\mathbf{I} := \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$ . However, our Bézout Lemma proof only used  $\mathbf{B}(0)$  and  $\mathbf{B}(1)$ , i.e that

$$*: \quad M_0 \cdot R_0 = R_0.$$

and so other values of  $M_0$  are possible.

As an example, the usual LBolt for  $\text{GCD}(3, 2)$  is

8a:

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	3	—	1	0
1	2	1	0	1
2	1	2	1	-1
3	0	$\infty$	-2	3

Another initial-matrix is  $M_0 := \begin{bmatrix} 3 & -3 \\ 0 & 1 \end{bmatrix}$ , yielding

8b:

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	3	—	3	-3
1	2	1	0	1
2	1	2	3	-4
3	0	$\infty$	-6	9

Row-2 gives us a *different* Bézout pair. We might conjecture that check-pair  $(-6, 9)$  equals the check-pair  $(-2, 3)$  from the first table, but multiplied by  $\text{Det}(M_0)$ .

Yet another init-matrix is  $M_0 := \begin{bmatrix} 7 & -9 \\ 2 & -2 \end{bmatrix}$ , producing

8c:

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	3	—	7	-9
1	2	1	2	-2
2	1	2	5	-7
3	0	$\infty$	-8	12

Row-2 gives us a *third* Bézout pair. The check-pair  $(-8, 12)$  indeed equals  $\text{Det}\left(\begin{smallmatrix} 7 & -9 \\ 2 & -2 \end{smallmatrix}\right)$  times the  $(-2, 3)$  from our first table.

In this last example

8d:

$n$	$r_n$	$q_n$	$s_n$	$t_n$
0	3	—	-5	9
1	2	1	-4	7
2	1	2	-1	2
3	0	$\infty$	-2	3

has  $M_0 := \begin{bmatrix} -5 & 9 \\ -4 & 7 \end{bmatrix}$ , whose determinant is 1, hence yielding the same check-tuple  $(-2, 3)$  as in table (8a).

**Check-row.** We study an LBolt seeded with a coprime pair  $r_0 \perp r_1$ , and initial-matrix  $\mathbf{M}_0$  st.  $(*)$  holds.

For  $k \geq 1$ , let

$$\mathbf{Q}_k := \begin{bmatrix} 0 & 1 \\ 1 & -\mathbf{q}_k \end{bmatrix}$$

and observe  $\text{Det}(\mathbf{Q}_k) = -1$ . Define product matrix

$$\mathbf{P}_n := \mathbf{Q}_n \cdots \mathbf{Q}_2 \mathbf{Q}_1;$$

hence  $\mathbf{P}_0$ , the empty product, is the identity matrix  $\mathbf{I}$ .

Update-rule (3) tells us that

$$\mathbf{R}_k = \mathbf{Q}_k \cdot \mathbf{R}_{k-1} \quad \text{and} \quad \mathbf{M}_k = \mathbf{Q}_k \cdot \mathbf{M}_{k-1}.$$

Consequently,

$$\begin{aligned} 9a: \quad \mathbf{R}_n &= \mathbf{P}_n \cdot \mathbf{R}_0, & \mathbf{M}_n &= \mathbf{P}_n \cdot \mathbf{M}_0 \\ &\text{and} \quad \text{Det}(\mathbf{M}_n) &= \text{Det}(\mathbf{M}_0) \cdot [-1]^n. \end{aligned}$$

Moreover,

$$9b: \quad \mathbf{M}_n \cdot \mathbf{R}_0 = \mathbf{P}_n \mathbf{M}_0 \cdot \mathbf{R}_0 \stackrel{\text{by } (*)}{=} \mathbf{P}_n \mathbf{R}_0 = \mathbf{R}_n.$$

Letting  $N := \text{GCD-index}$ , we have that

$$**: \quad \mathbf{M}_N \cdot \mathbf{R}_0 = \mathbf{R}_N \stackrel{\text{recall}}{=} \begin{bmatrix} 1 \\ 0 \end{bmatrix},$$

since  $r_0 \perp r_1$ . Have  $(S, T) := (s_N, t_N)$  denote the Bézout-pair, and let  $(\alpha, \beta) := (s_{N+1}, t_{N+1})$  be the pair whose values we wish to determine. Finally, set  $\delta := \text{Det}(\mathbf{M}_0)$ .

Our  $(**)$  gives the top two lines of

$$\begin{aligned} Sr_0 + Tr_1 &= 1, \\ \alpha r_0 + \beta r_1 &= 0. \quad \text{Notice that} \\ -\alpha T + \beta S &= \delta \cdot [-1]^N \end{aligned}$$

courtesy (9a), since  $\text{Det}(\mathbf{M}_N) = \text{Det}\left(\begin{bmatrix} S & T \\ \alpha & \beta \end{bmatrix}\right)$ . Multiplying the middle eqn by  $T$  and the bottom by  $r_0$  gives

$$\begin{aligned} \alpha Tr_0 + \beta Tr_1 &= 0 \quad \text{and} \\ -\alpha Tr_0 + \beta Sr_0 &= r_0 \delta [-1]^N. \end{aligned}$$

Adding them yields

$$\beta \stackrel{\text{note}}{=} \beta \cdot [Sr_0 + Tr_1] = r_0 \delta \cdot [-1]^N.$$

Finally, plugging this into the middle eqn gives

$$0 = \alpha r_0 + r_0 \delta [-1]^N \cdot r_1.$$

When  $r_0 \neq 0$ , then  $0 = \alpha + r_1 \delta [-1]^N$ . Hence

$$\alpha = -r_1 \delta \cdot [-1]^N = r_1 \delta \cdot [-1]^{N+1}.$$

We have proven the following theorem.

**9c: Check-value Theorem.** Consider an LBolt seeded with integers  $r_0 \neq 0$  and  $r_1$ , together with an initial-matrix  $\mathbf{M}_0$  satisfying

$$9d: \quad \mathbf{M}_0 \cdot \begin{bmatrix} r_0 \\ r_1 \end{bmatrix} = \begin{bmatrix} r_0 \\ r_1 \end{bmatrix}.$$

Let  $N := \text{GCD-index}$  and  $\mathcal{G} := \text{GCD}(r_0, r_1)$ . Then

$$\begin{aligned} 9e: \quad s_{N+1} \cdot \mathcal{G} &= r_1 \cdot \text{Det}(\mathbf{M}_0) \cdot [-1]^{N+1} \\ \text{and } t_{N+1} \cdot \mathcal{G} &= r_0 \cdot \text{Det}(\mathbf{M}_0) \cdot [-1]^N. \end{aligned}$$

(Recall that our standard LBolt has  $\text{Det}(\mathbf{M}_0) = 1$ .)  $\diamond$

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